

# Exact schedulability test for sporadic mixed-criticality real-time systems using antichains and oracles

**Simon Picard**, Antonio Paolillo, Gilles Geeraerts, Joël Goossens

Ourway, Vrije Universiteit Brussel, Université libre de Bruxelles

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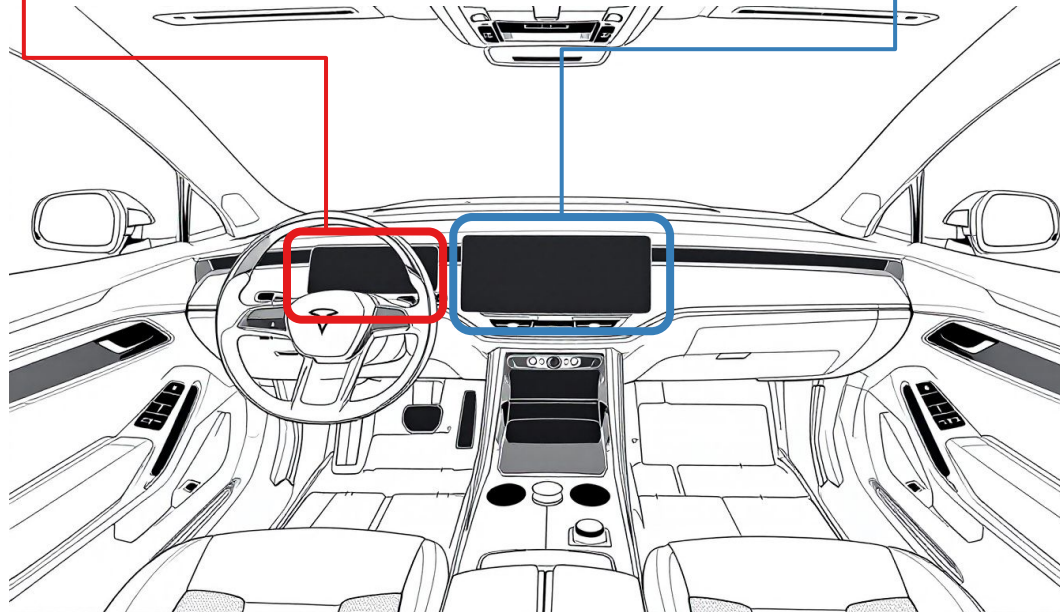


67 10 96 5

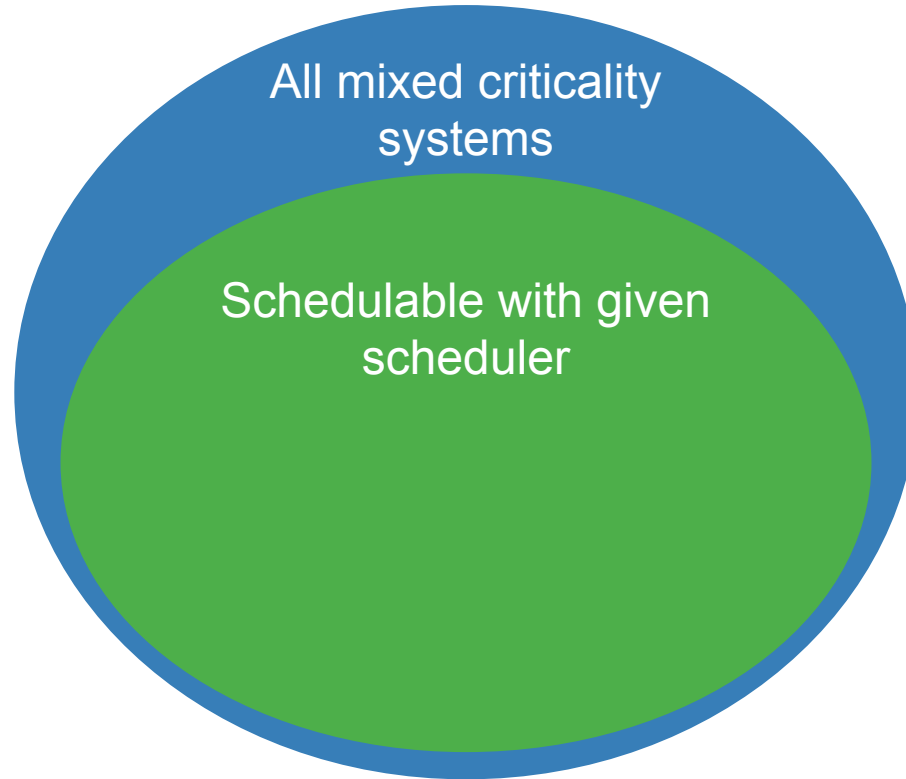
# Introduction to mixed-criticality systems

automated driving

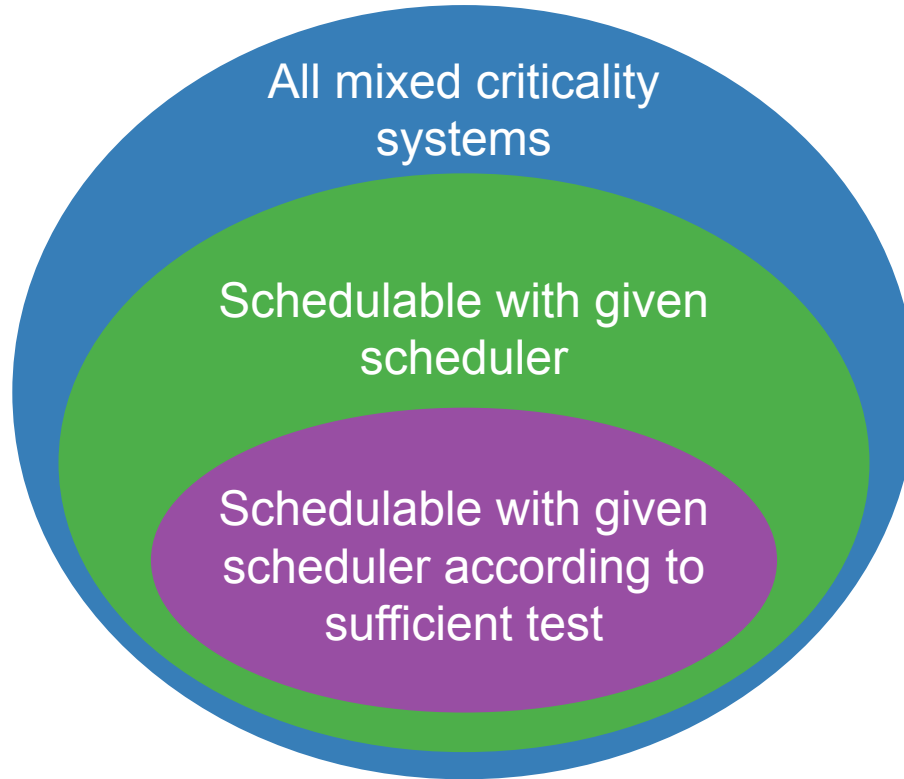
passenger entertainment



# Problem statement



# Problem statement



# Sporadic dual-criticality systems

Task set  $\tau = \{\tau_1, \tau_2, \dots, \tau_n\}$

Task  $\tau_i = \langle\langle L_i, T_i, D_i, C_i(\text{LO}), C_i(\text{HI}) \rangle\rangle$

$L_i$  : criticality level (**LO** or **HI**)

$T_i$  : minimum interarrival time (**period**) between jobs released **sporadically**

$D_i$  : relative **deadline** (constrained)

$C_i$  : worst-case execution time (WCET) **per criticality level**

Single processor

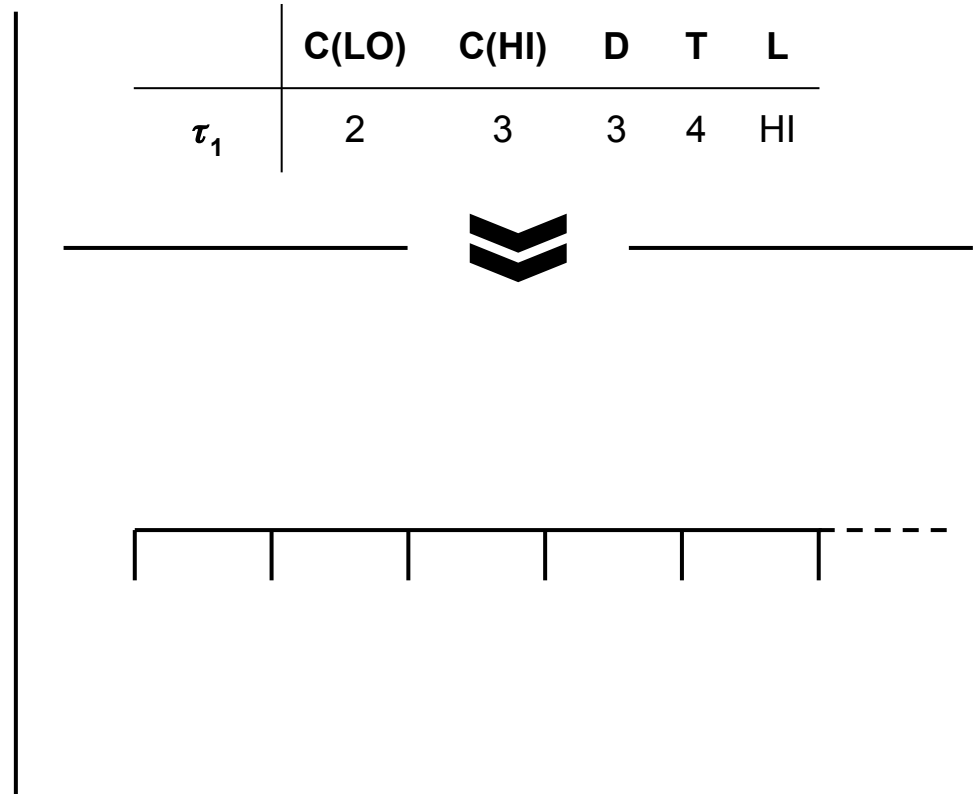
# Lifecycle of a sporadic dual-criticality system

At each **clock tick**:

1. **Potential** job release
2. Scheduling
3. Execution
4. **Potential** completion signal
5. Possible criticality mode change

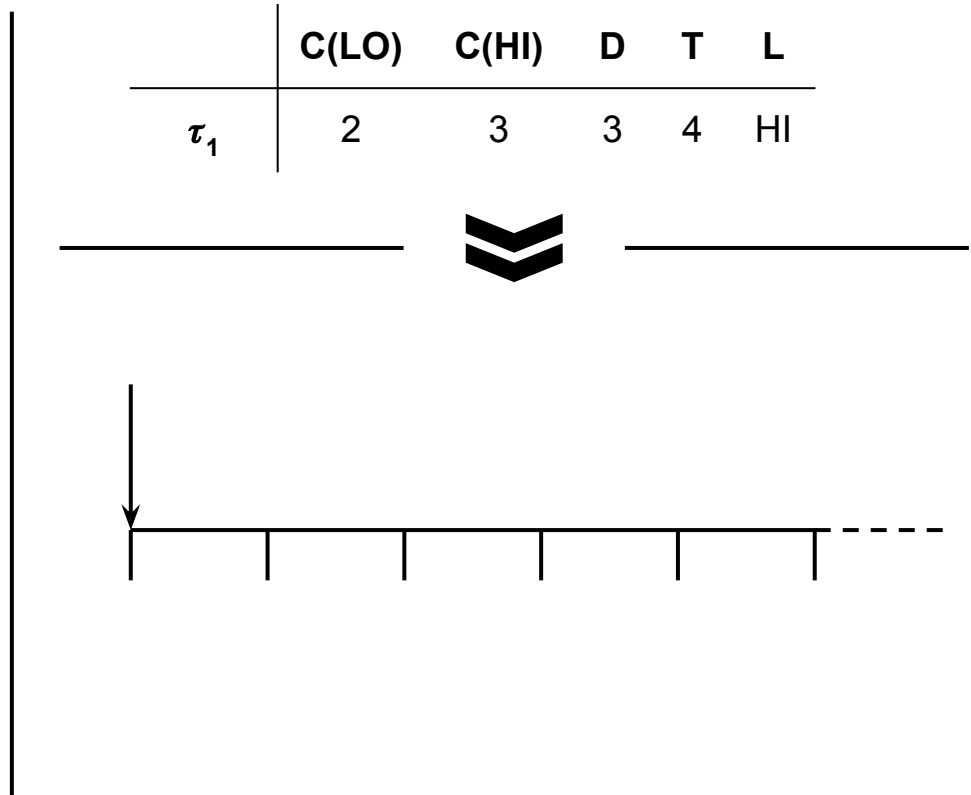
**Non-deterministic job release and completion signal → multiple scenarios**

# Sporadic dual-criticality task chronogram



# Sporadic dual-criticality task chronogram

↓ : release

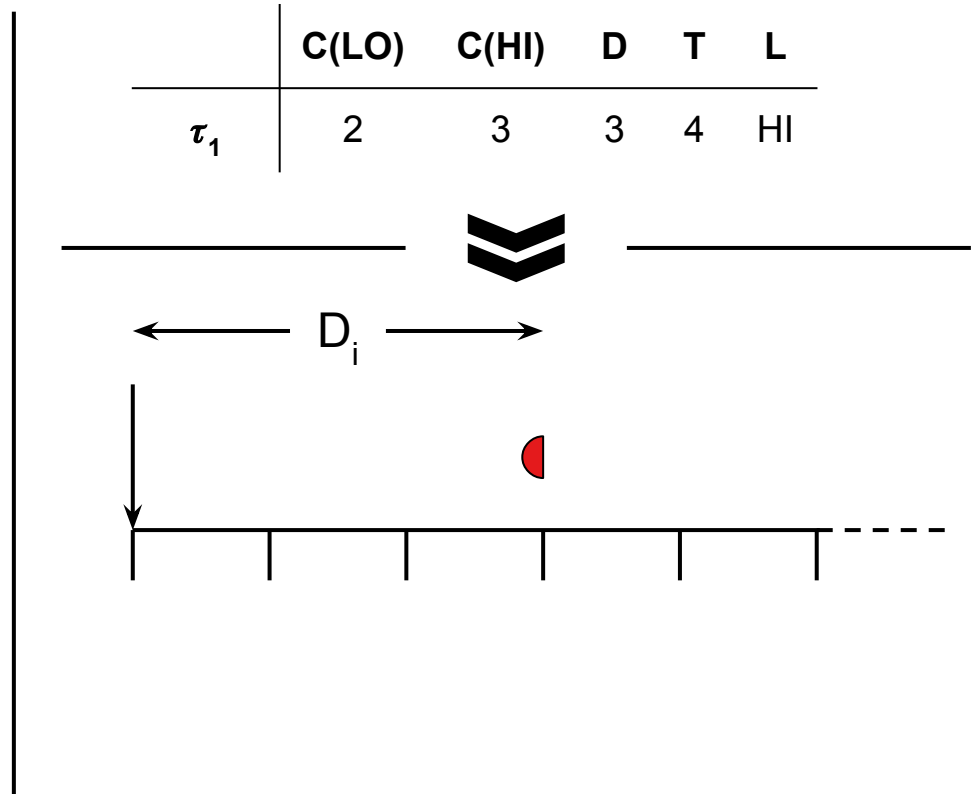




# Sporadic dual-criticality task chronogram

↓: release

◐: absolute deadline

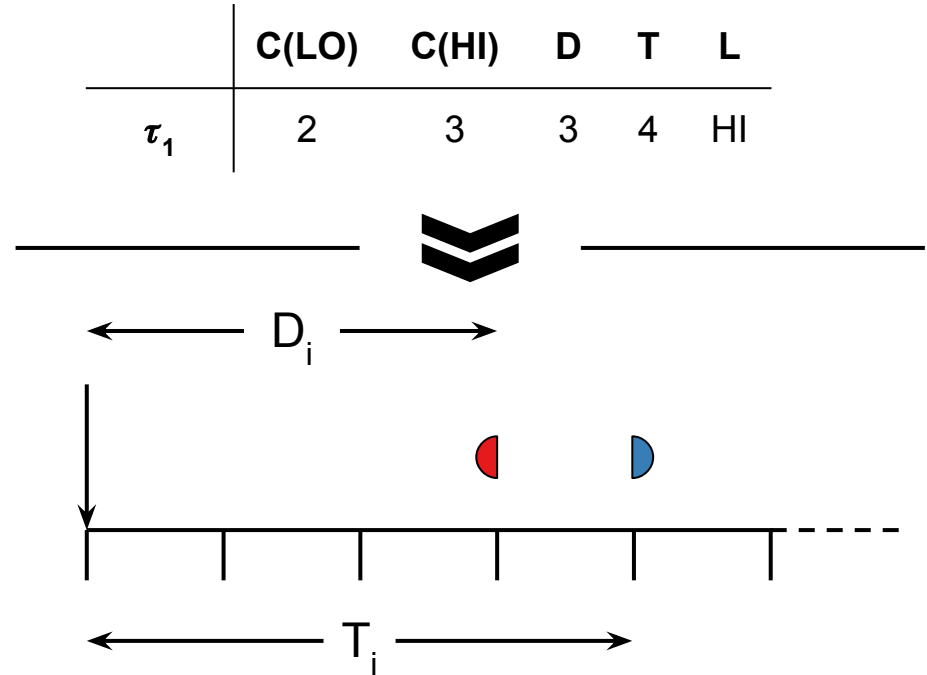


# Sporadic dual-criticality task chronogram

↓: release

◐: absolute deadline

◑: earliest next arrival time



# Sporadic dual-criticality task chronogram

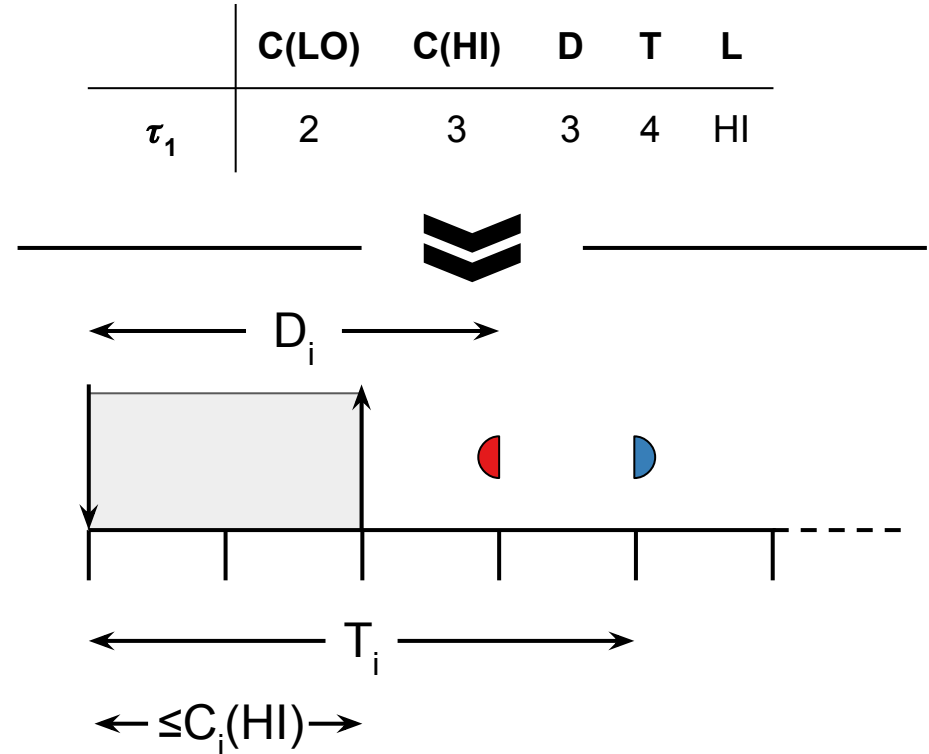
↓: release

◐: absolute deadline

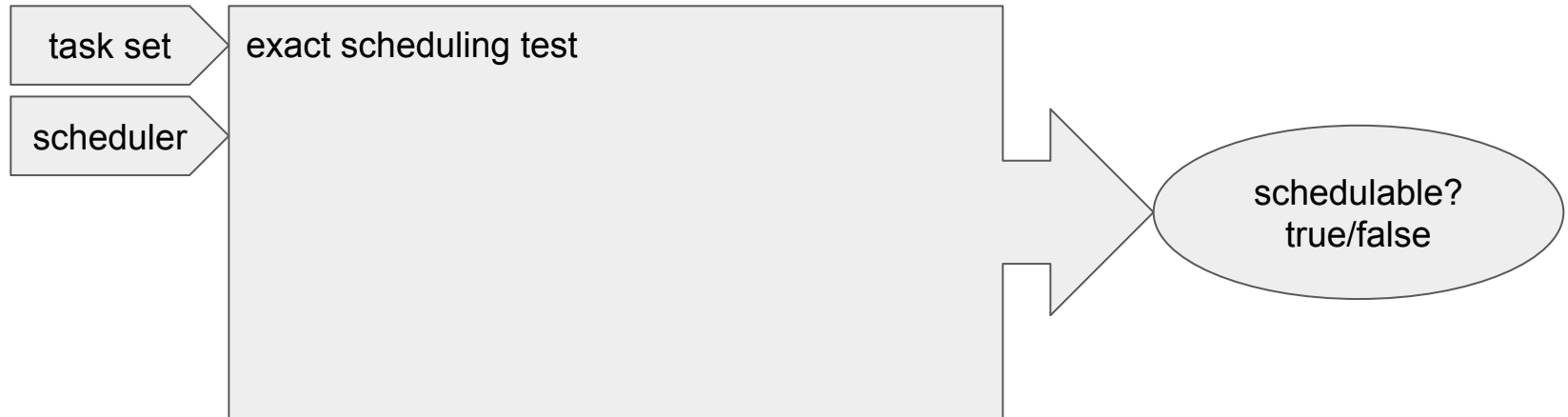
◑: earliest next arrival time

◻: execution

↑: completion signal



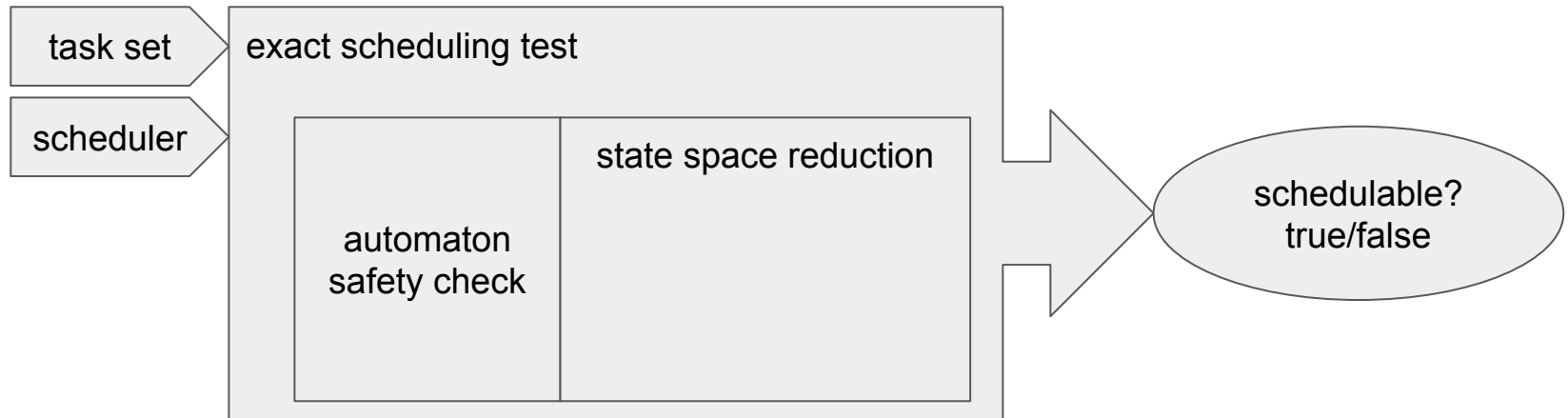
# Objective of the work



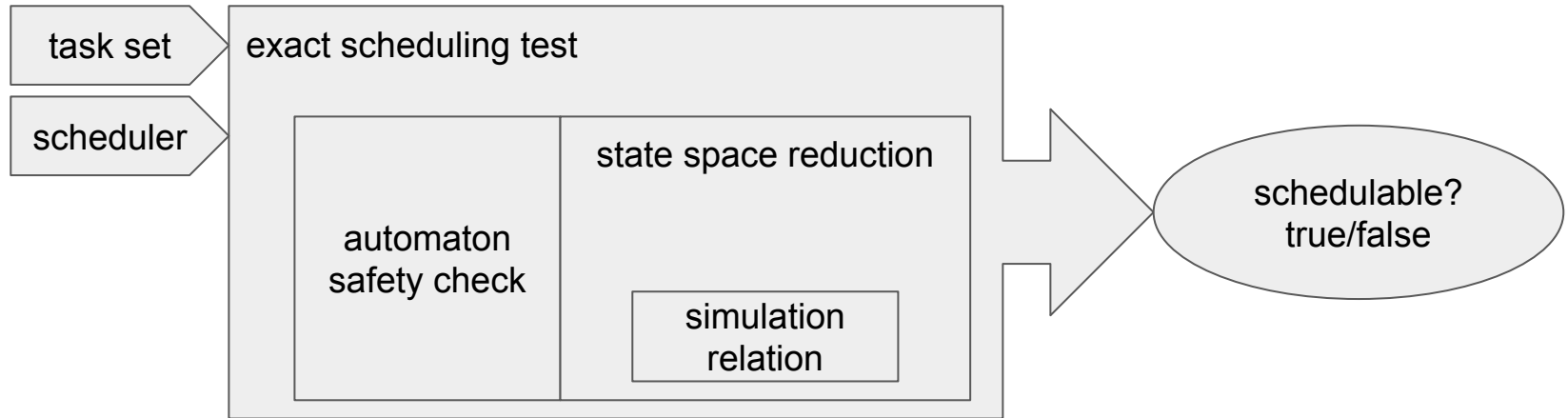
# Objective of the work



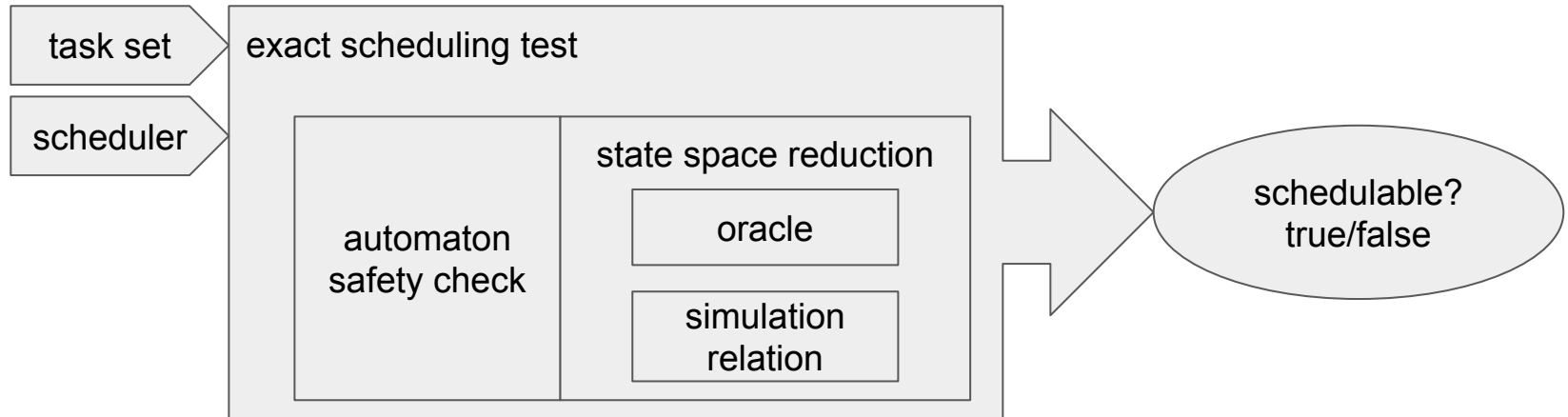
# Objective of the work



# Objective of the work

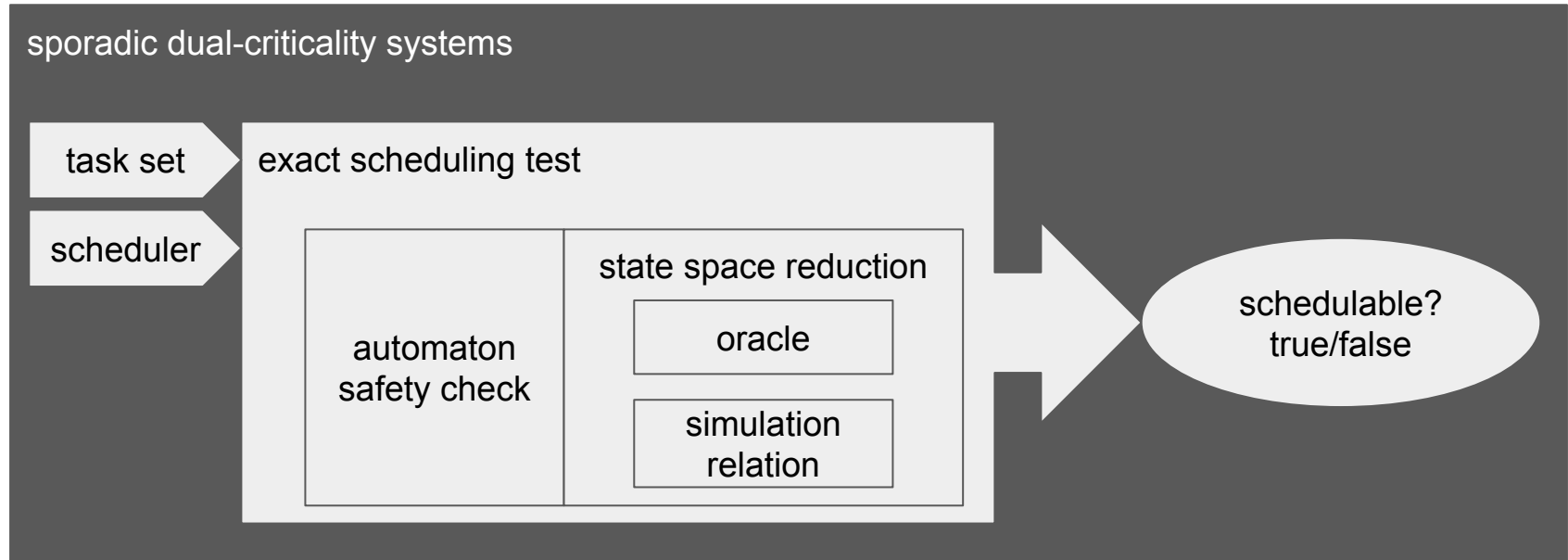


# Objective of the work

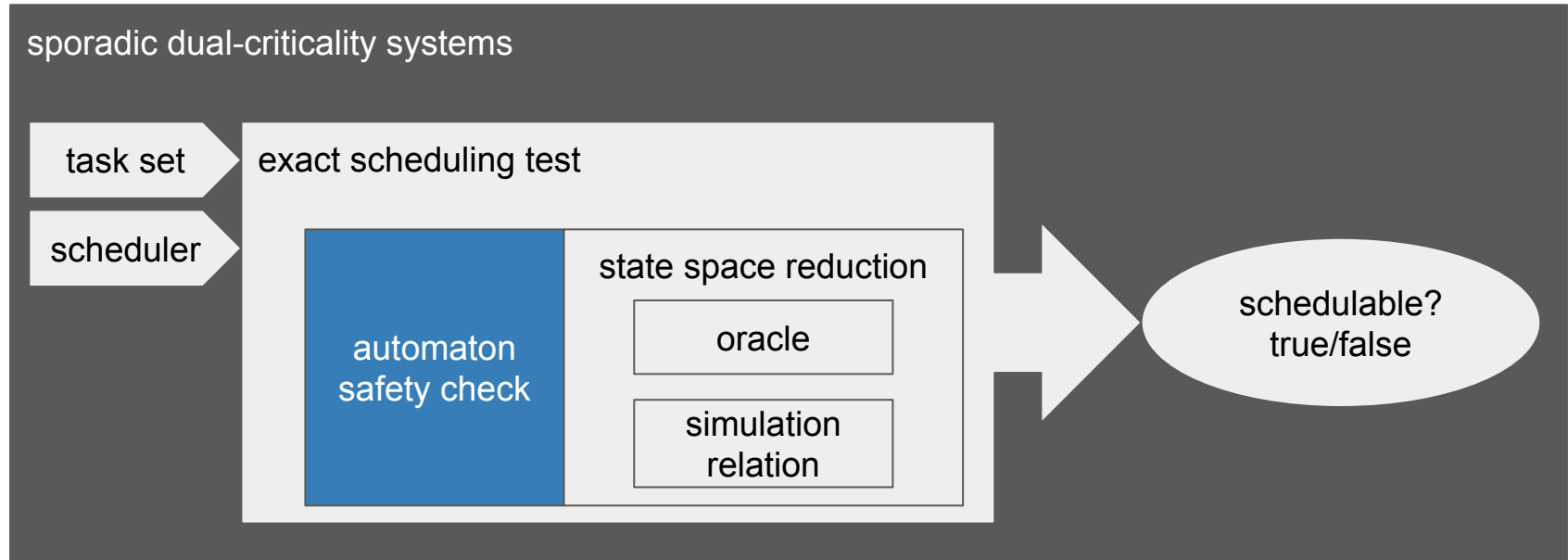




# Objective of the work

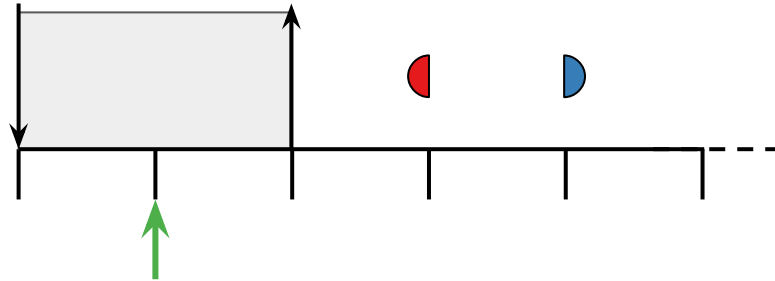


# Objective of the work



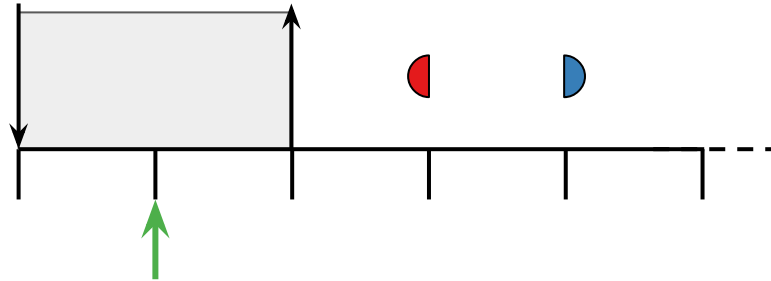
# States of the automaton

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



# Criticality (cri)

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI

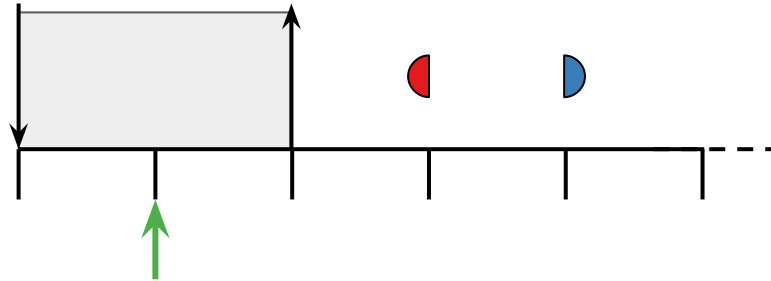


cri = HI if a mode change happened, LO otherwise

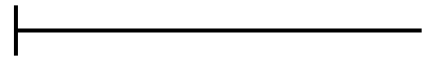
cri = LO

# Next Arrival Time (nat)

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



nat =

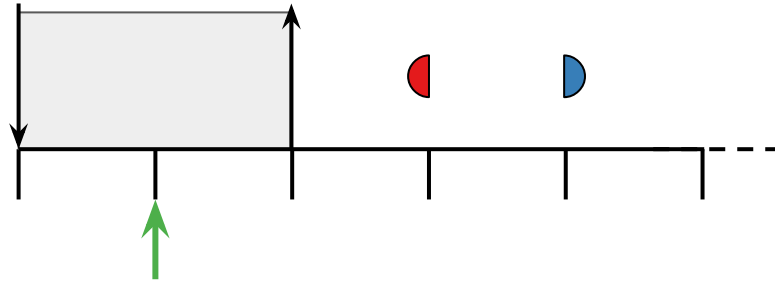


nat =

3

# Remaining Computing Time (rct) for the current criticality

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI

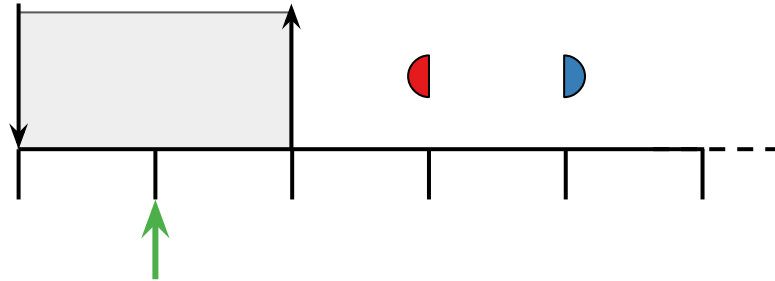


$$\text{rct} = - \text{---} + C(\text{cri})$$

$$\text{rct} = - 1 + 2 = 1$$

# States of the automaton

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI

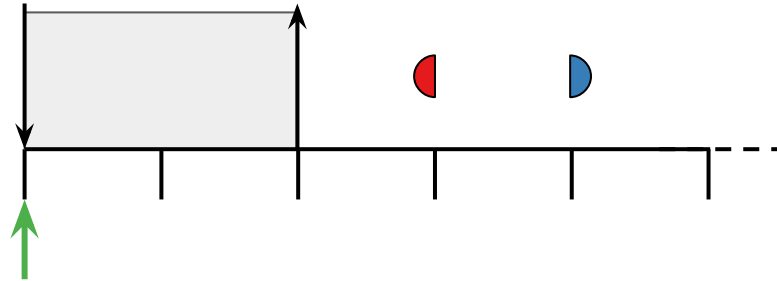


cri (rct<sub>i</sub>, nat<sub>i</sub>)

LO (1, 3)

# From chronogram to path

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI

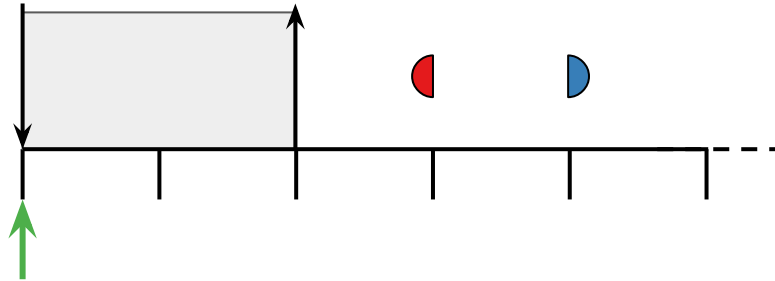


cri  
(rct<sub>1</sub>, nat<sub>1</sub>)



# From chronogram to path

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI

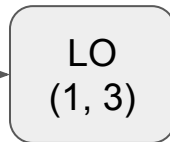
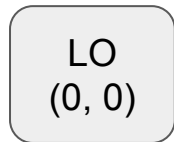
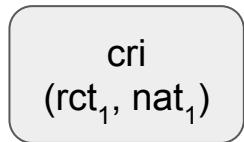
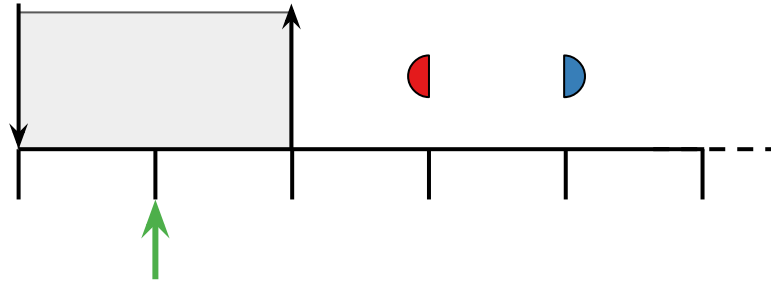


cri  
( $rct_1$ ,  $nat_1$ )

LO  
(0, 0)

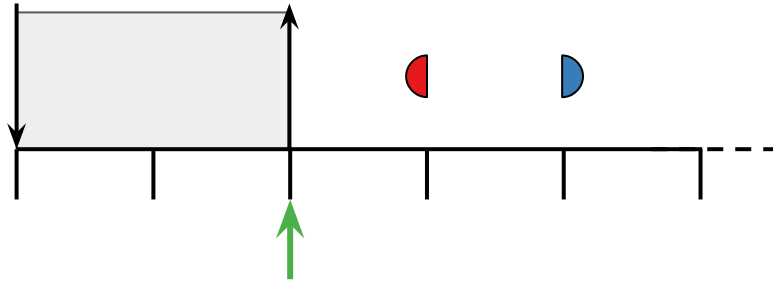
# From chronogram to path

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



# From chronogram to path

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



cri  
( $rct_1$ ,  $nat_1$ )

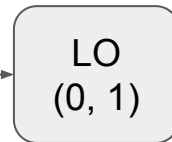
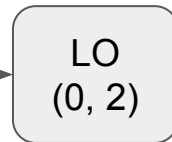
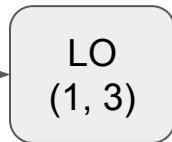
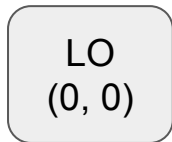
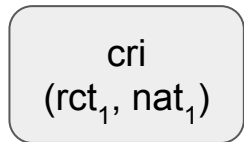
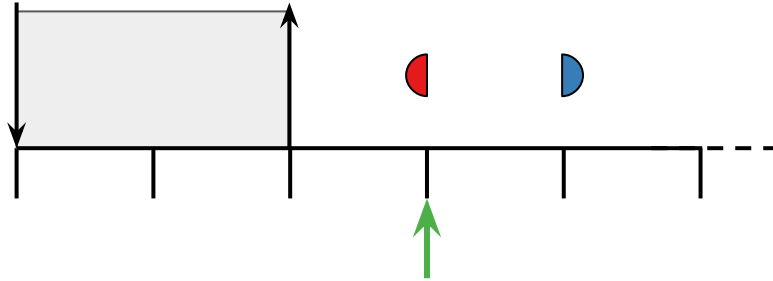
LO  
(0, 0)

LO  
(1, 3)

LO  
(0, 2)

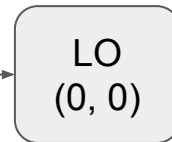
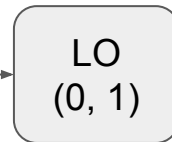
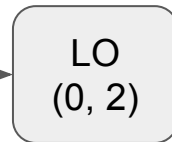
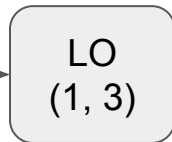
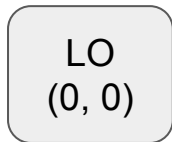
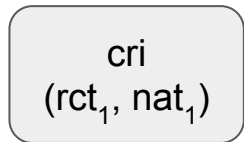
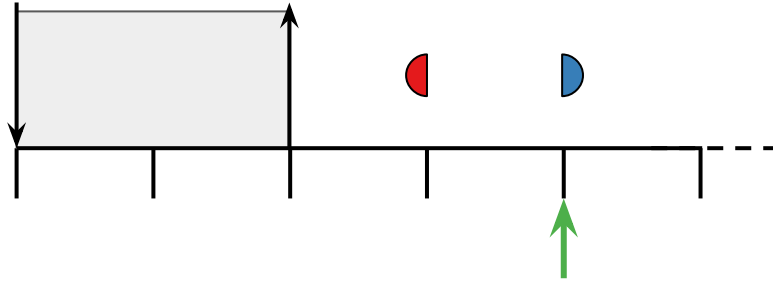
# From chronogram to path

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



# From chronogram to path

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



# From task set and scheduler to automaton

	$\tau_1$
<b>C(LO)</b>	2
<b>C(HI)</b>	3
<b>D</b>	3
<b>T</b>	4
<b>L</b>	HI

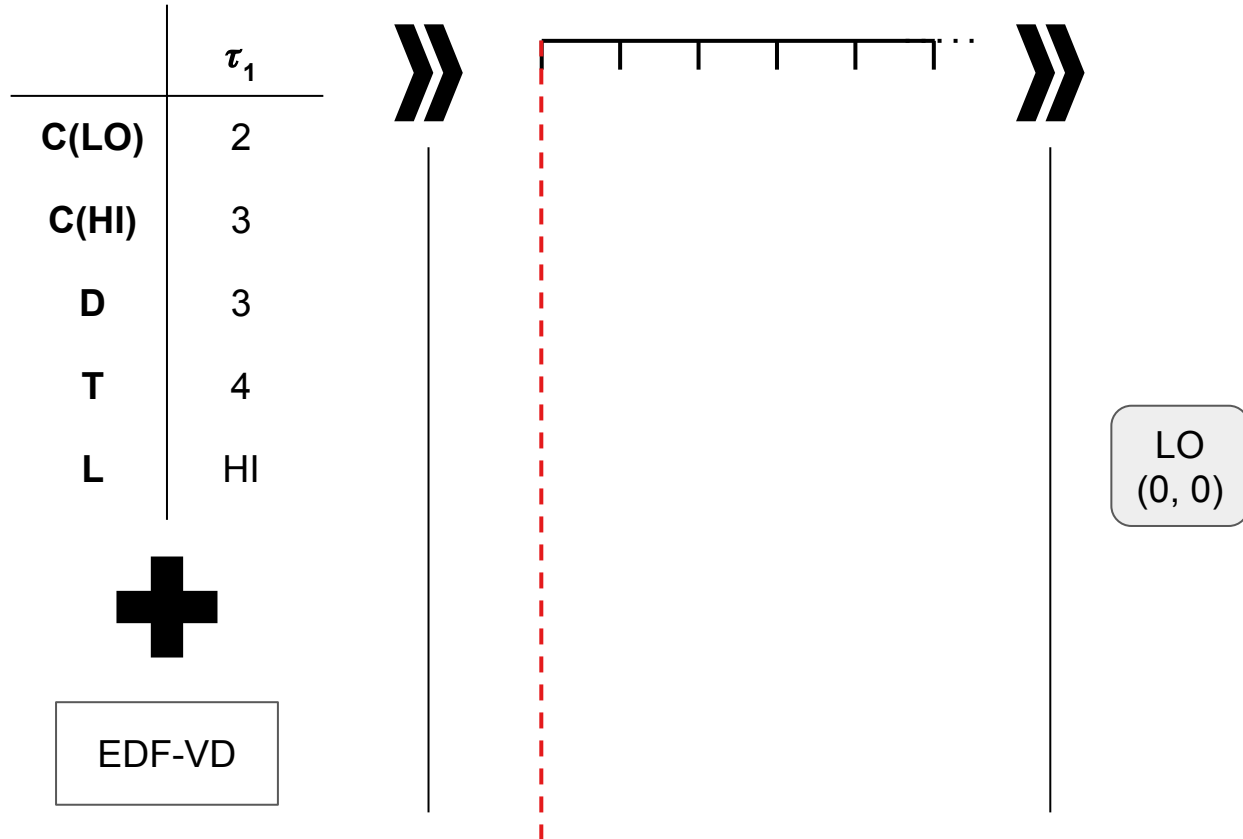
# From task set and scheduler to automaton

	$\tau_1$
<b>C(LO)</b>	2
<b>C(HI)</b>	3
<b>D</b>	3
<b>T</b>	4
<b>L</b>	HI



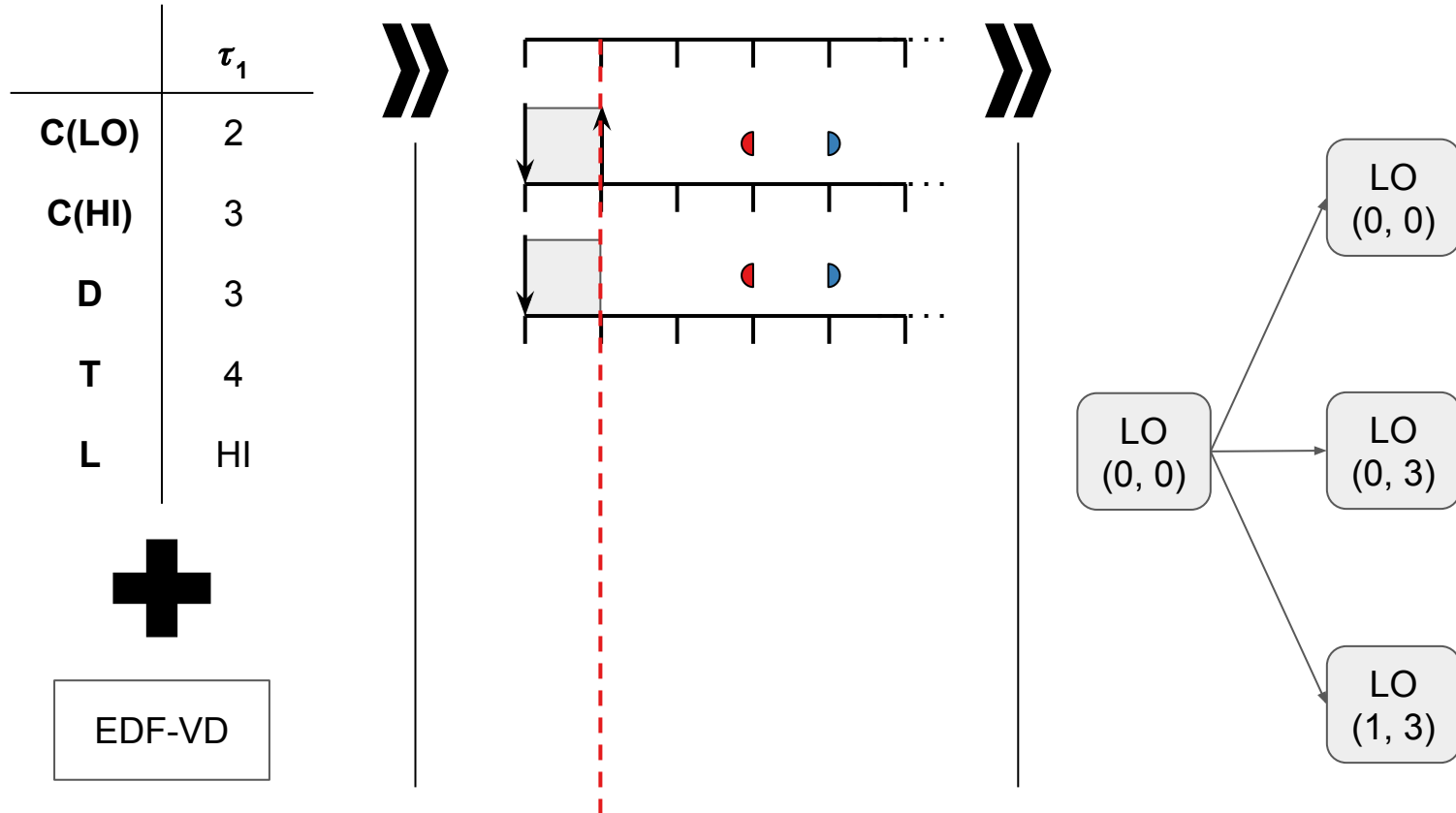
EDF-VD

# From task set and scheduler to automaton

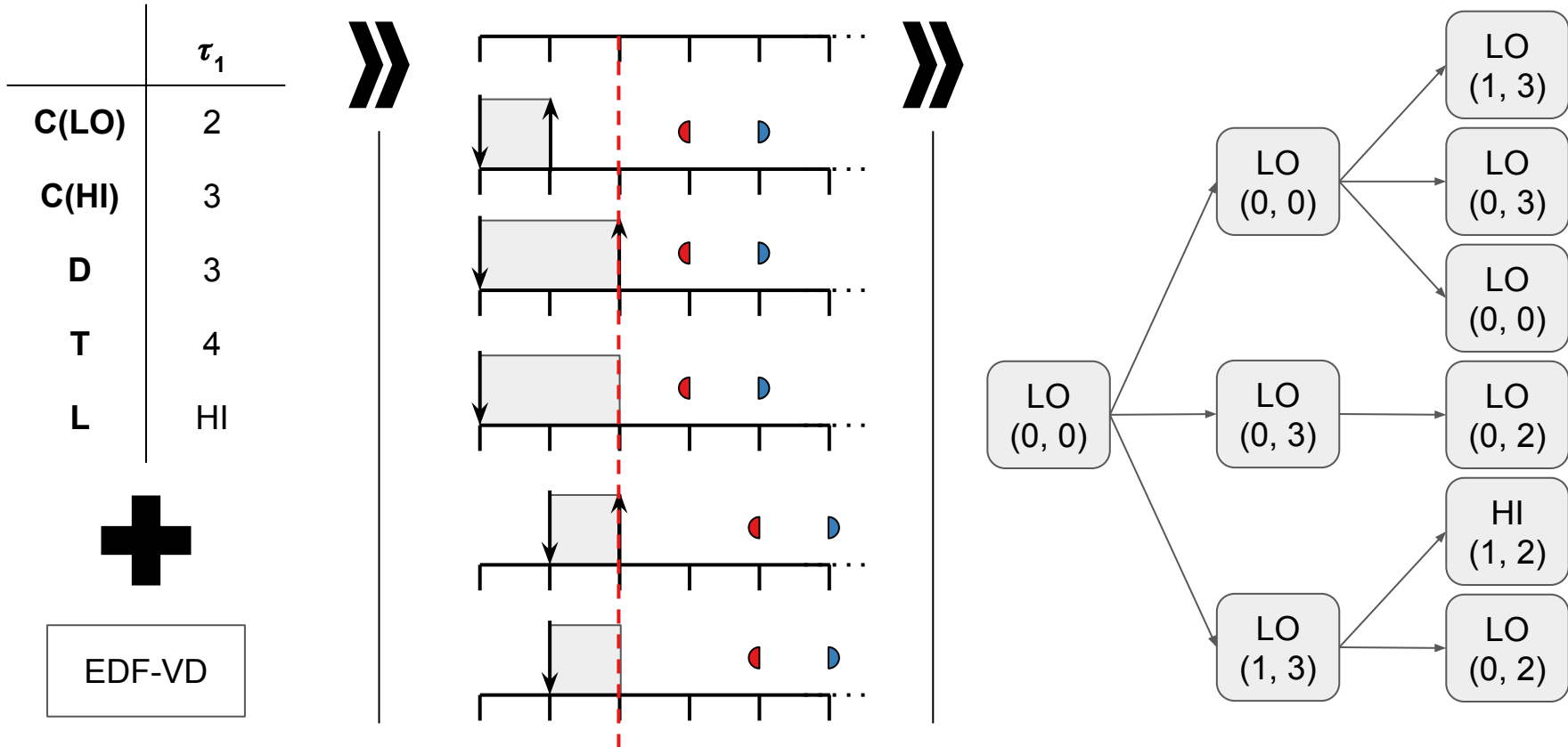




# From task set and scheduler to automaton



# From task set and scheduler to automaton

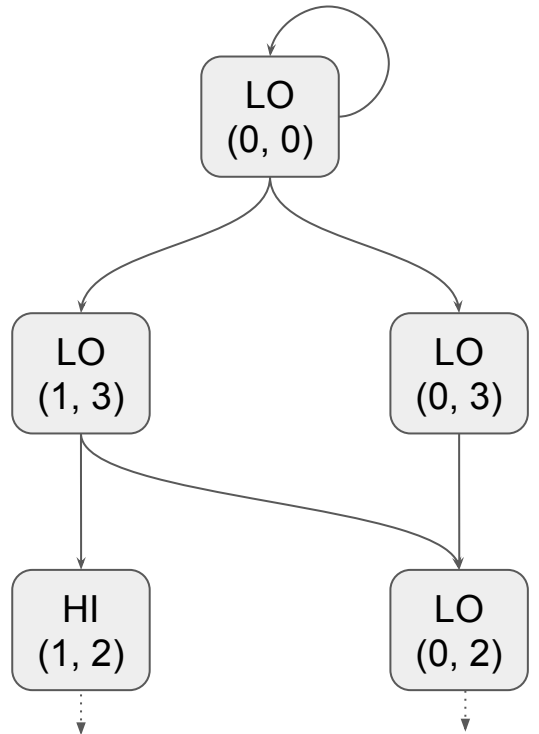
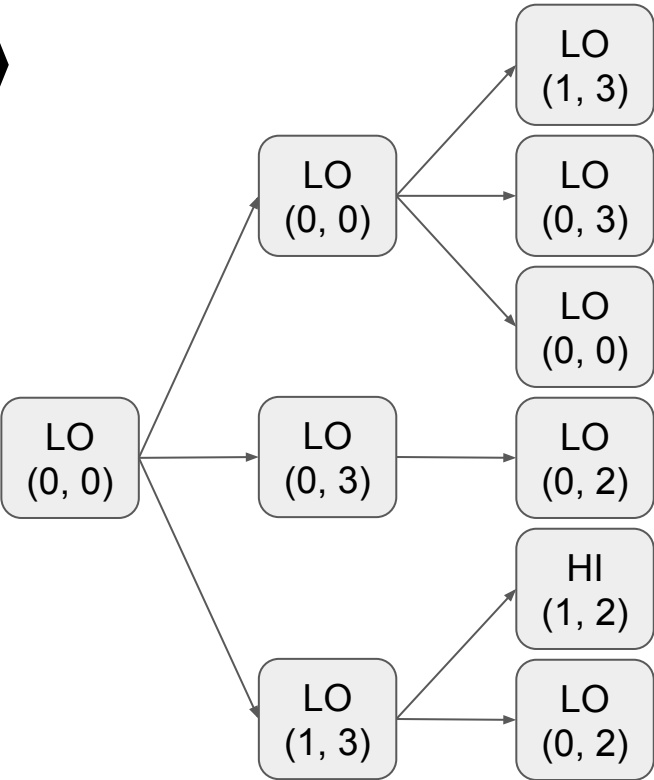


# From task set and scheduler to automaton

	$\tau_1$
C(LO)	2
C(HI)	3
D	3
T	4
L	HI



EDF-VD

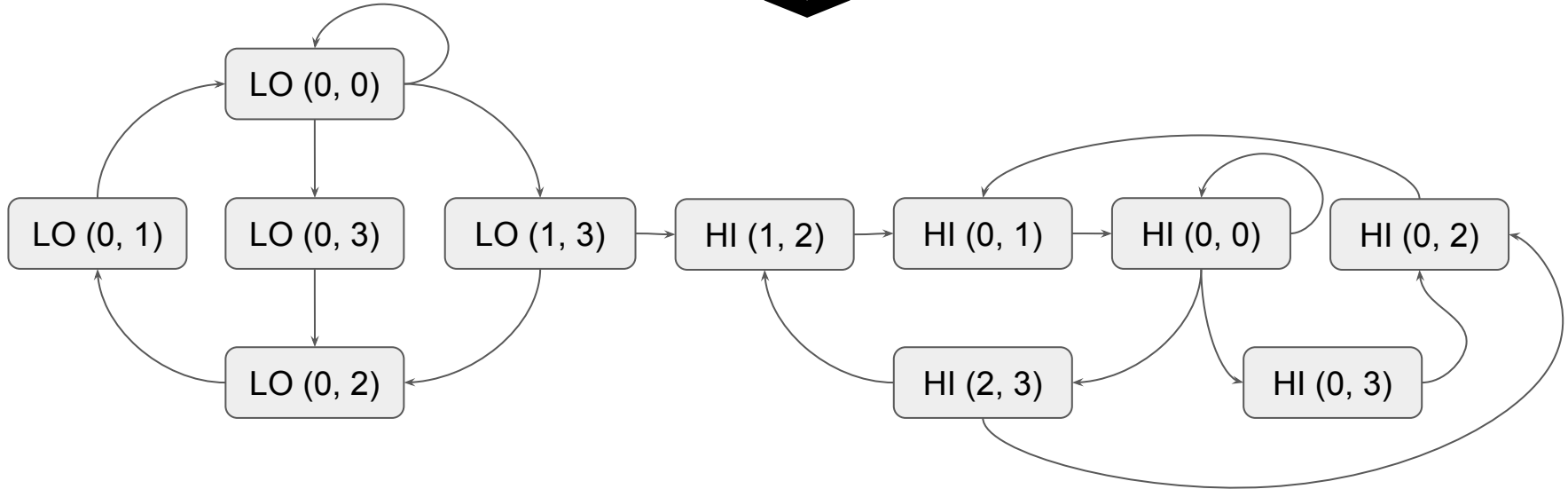


# From task set and scheduler to automaton

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



EDF-VD

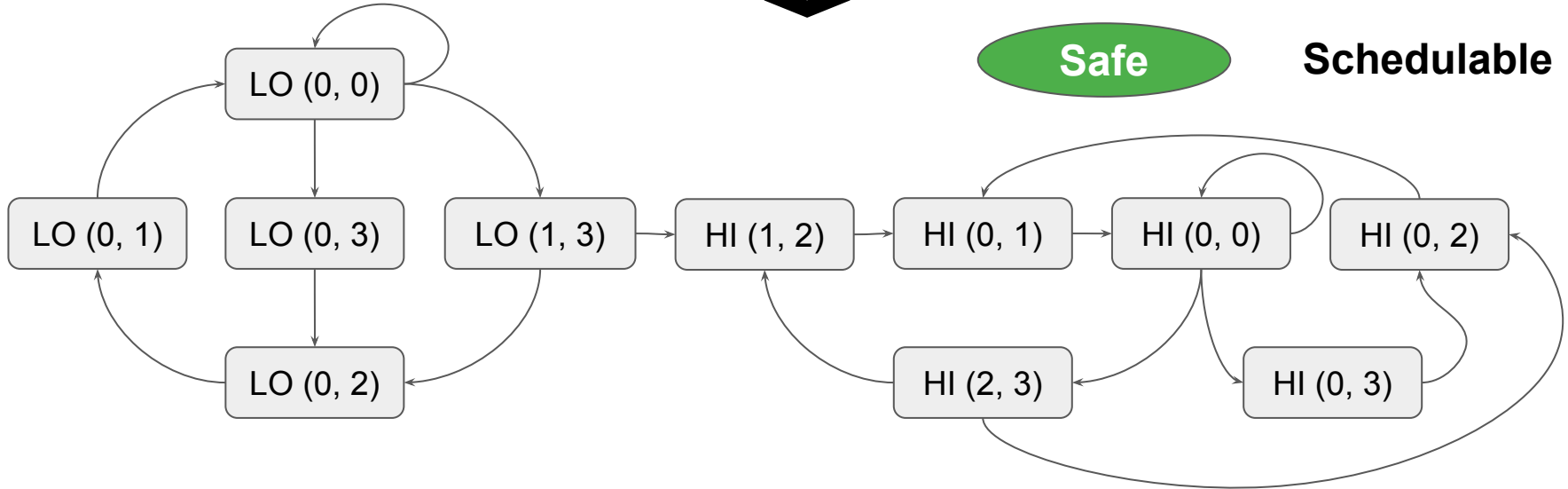


# From task set and scheduler to automaton

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



EDF-VD



# Unschedulable task set

	$\tau_1$	$\tau_2$
<b>C(LO)</b>	1	2
<b>C(HI)</b>	6	2
<b>D</b>	6	3
<b>T</b>	6	3
<b>L</b>	HI	LO



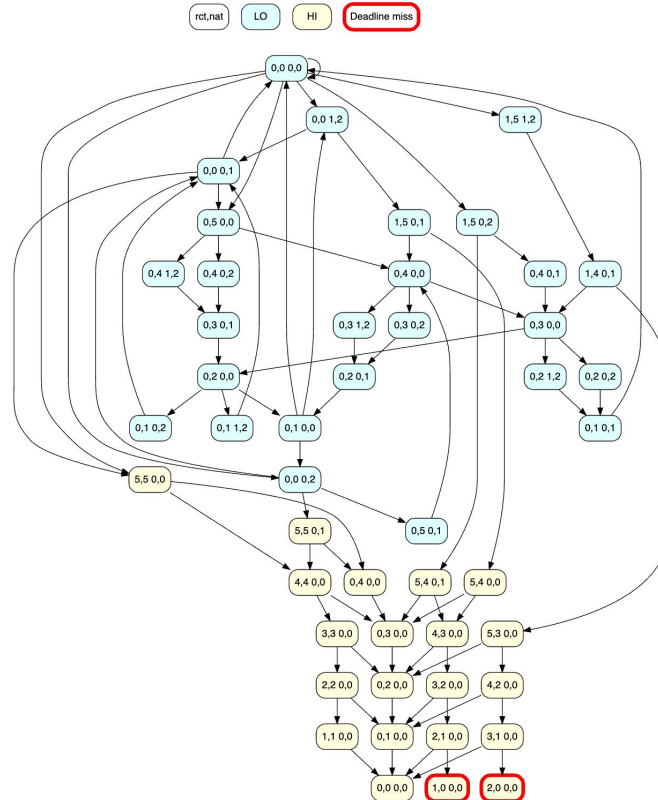
EDF-VD

# Unschedulable task set

	$\tau_1$	$\tau_2$
<b>C(LO)</b>	1	2
<b>C(HI)</b>	6	2
<b>D</b>	6	3
<b>T</b>	6	3
<b>L</b>	HI	LO



EDF-VD



**Unsafe**

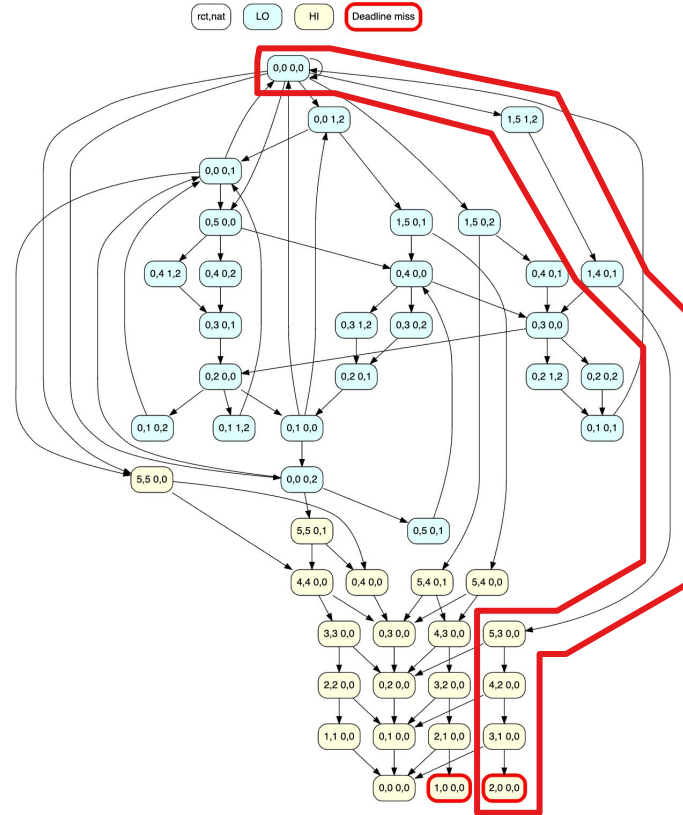
**Unschedulable**

# Unschedulable task set

	$\tau_1$	$\tau_2$
C(LO)	1	2
C(HI)	6	2
D	6	3
T	6	3
L	HI	LO



EDF-VD



Unsafe

Unschedulable



# Encountering a deadline miss (with EDF-VD)



	C(LO)	C(HI)	D	T	L
$\tau_1$	1	6	6	6	HI
$\tau_2$	2	2	3	3	LO



EDF-VD

- Both tasks release their job
- No early completion
- EDF-VD schedules  $\tau_2$  then  $\tau_1$  ( $\lambda=0.5$ )

# Encountering a deadline miss (with EDF-VD)



**at deadline**

time to deadline (ttd) = nat-T+D

$$ttd = 0-6+6 = 0$$



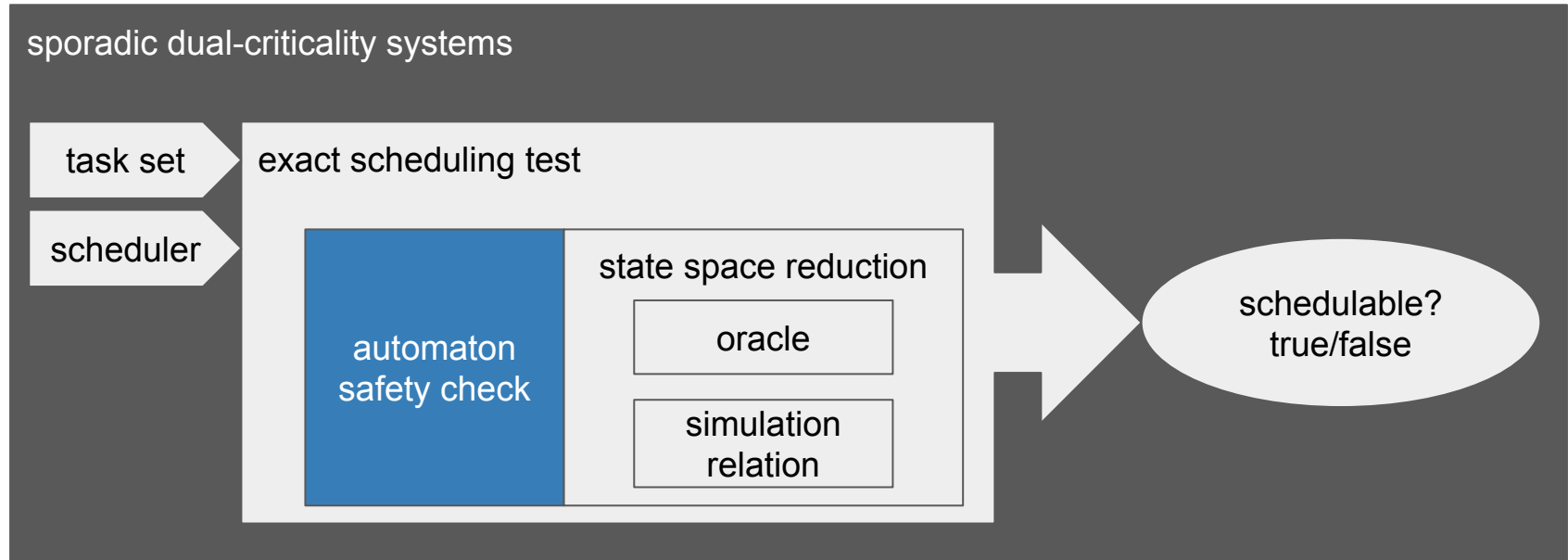
**not fully executed**

$$rct > 0$$

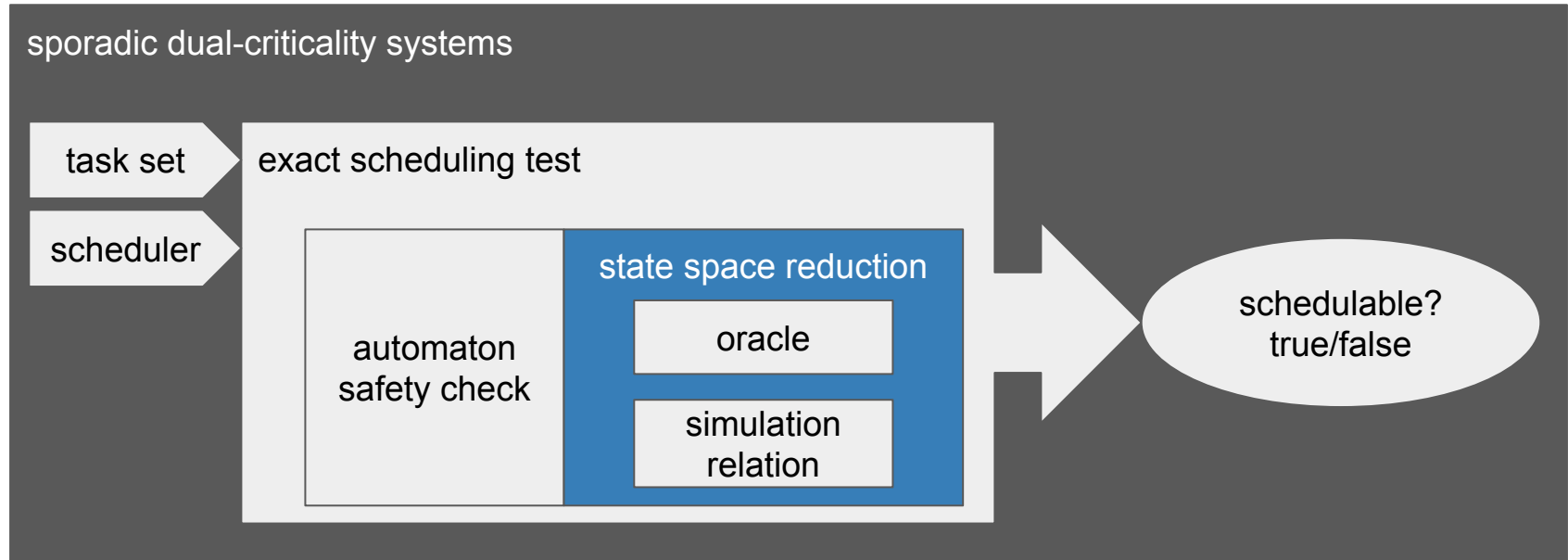


deadline miss

# Objective of the work



# Objective of the work



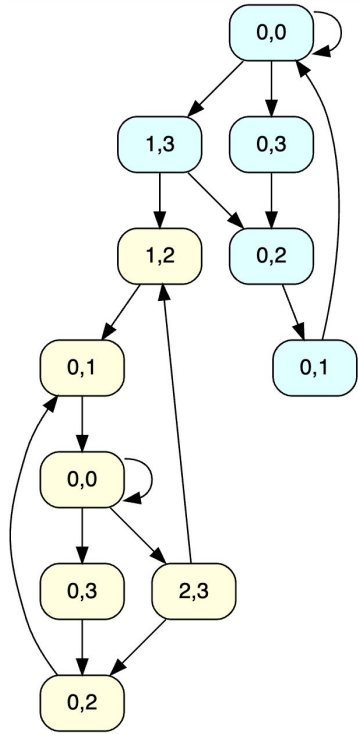
# State space exponential on the number of tasks

## 1 task

	$\tau_1$
C(LO)	2
C(HI)	3
D	3
T	4
L	HI



EDF-VD

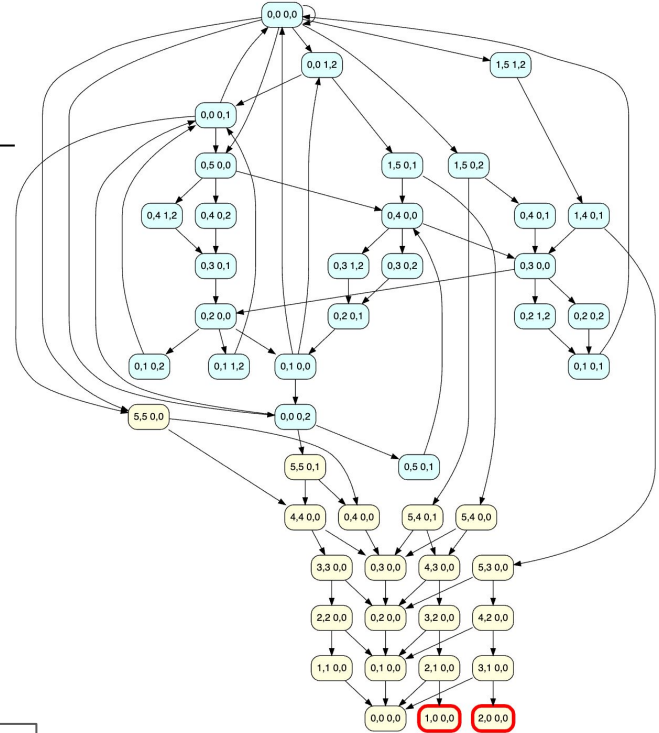


## 2 tasks

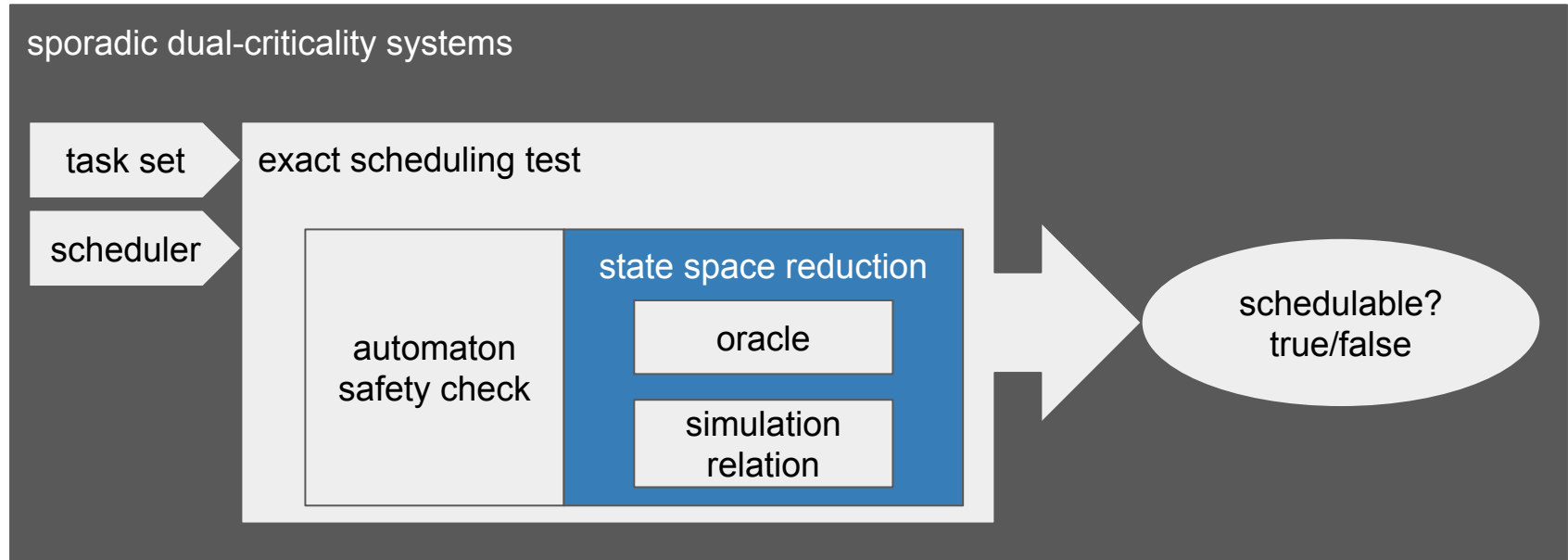
	$\tau_1$	$\tau_2$
C(LO)	1	2
C(HI)	6	2
D	6	3
T	6	3
L	HI	LO



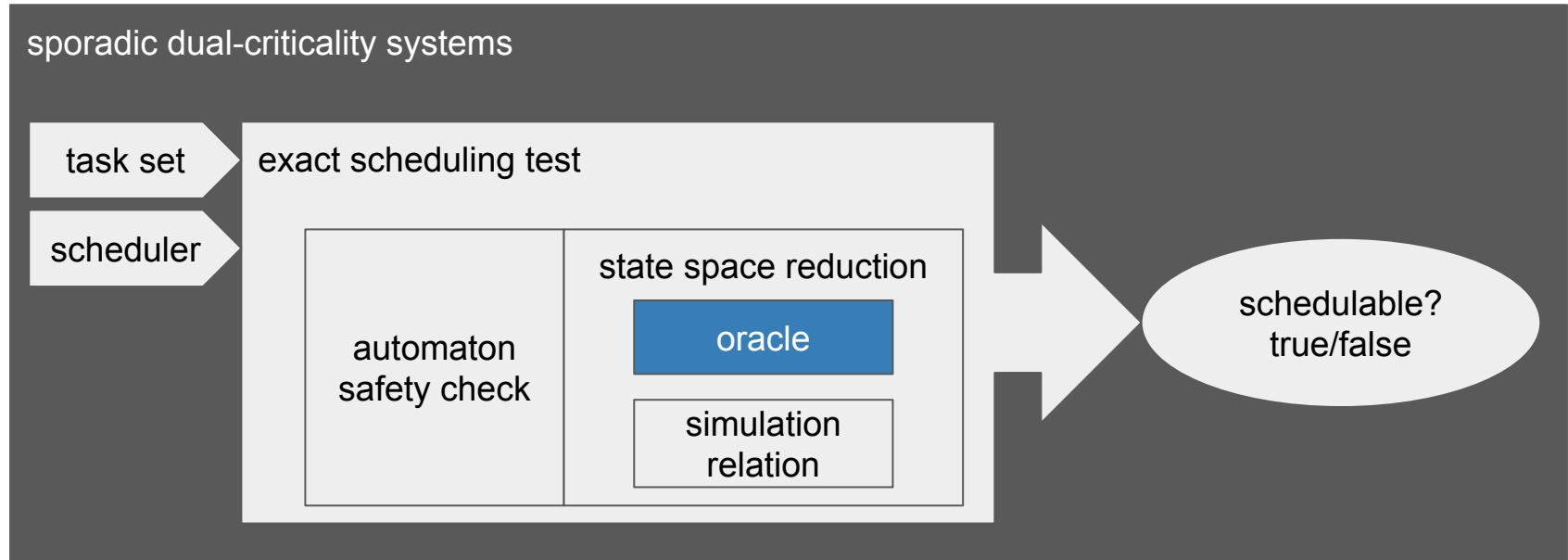
EDF-VD



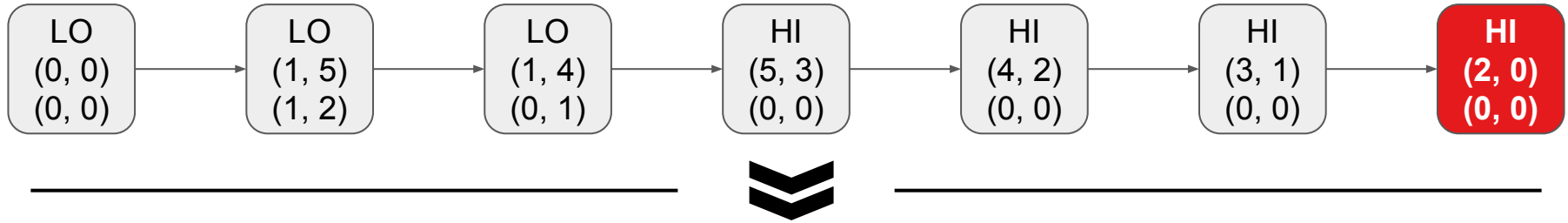
# Objective of the work



# Objective of the work



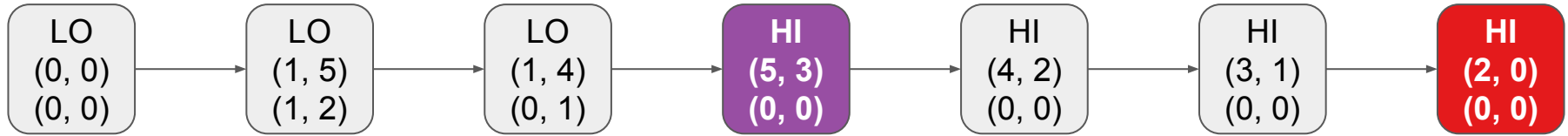
# Encountering a deadline miss (with EDF-VD)



Can an **oracle predict** a future deadline miss?



# Necessary condition to predict deadline miss (oracle)



**necessary condition**

$\text{laxity} \geq 0$

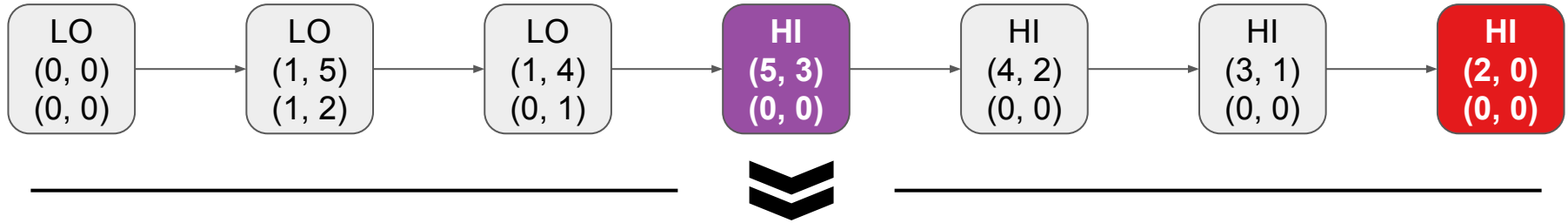
$\text{laxity} = \text{ttd} - \text{rct}$

$\text{laxity} = 3 - 5$

$\text{laxity} = -2 < 0$

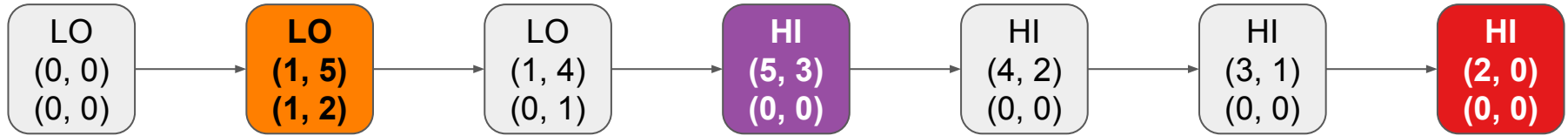
(future) deadline miss

# Predicting the deadline miss earlier?



Can another **oracle predict** a future deadline miss **even earlier**?

# Anticipating a future mode change for a stronger oracle



## necessary condition

worst laxity  $\geq 0$

worst laxity = laxity - (C(HI)-C(LO))

worst laxity = 4 - (6-1) < 0

worst laxity = 4 - 5 < 0

worst laxity = -1 < 0

(future) deadline miss

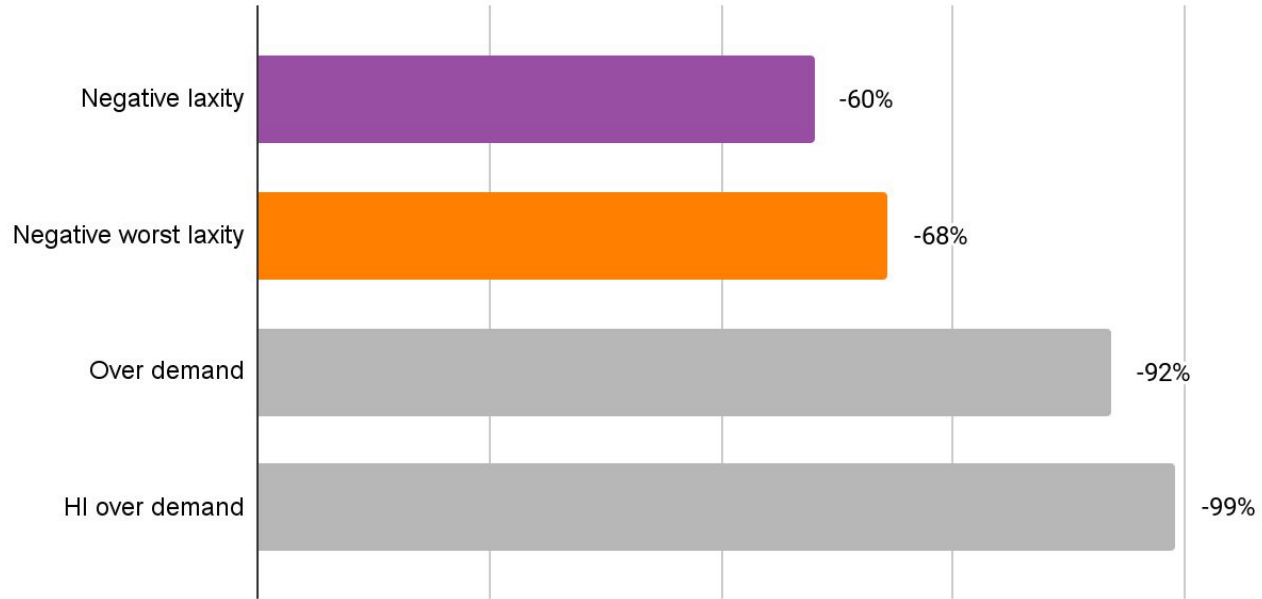
# Oracles' impact on unschedulable task sets

## Visited state reduction (median)

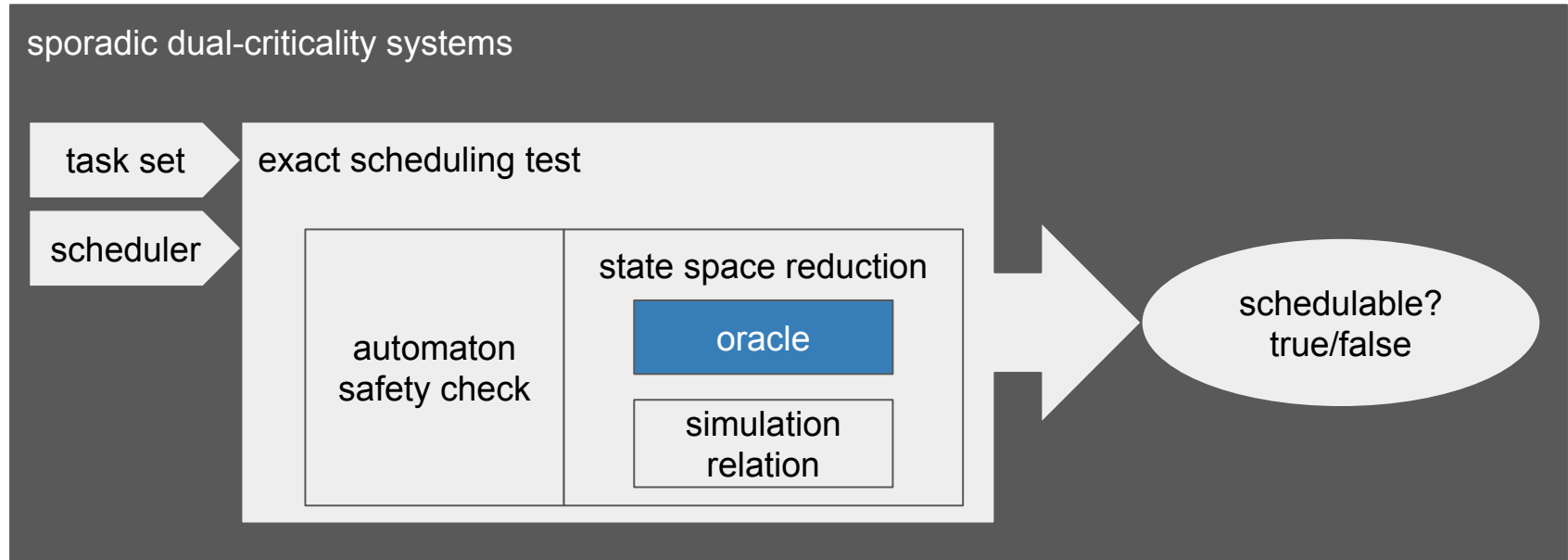


# Paper only oracles' impact on unschedulable task sets

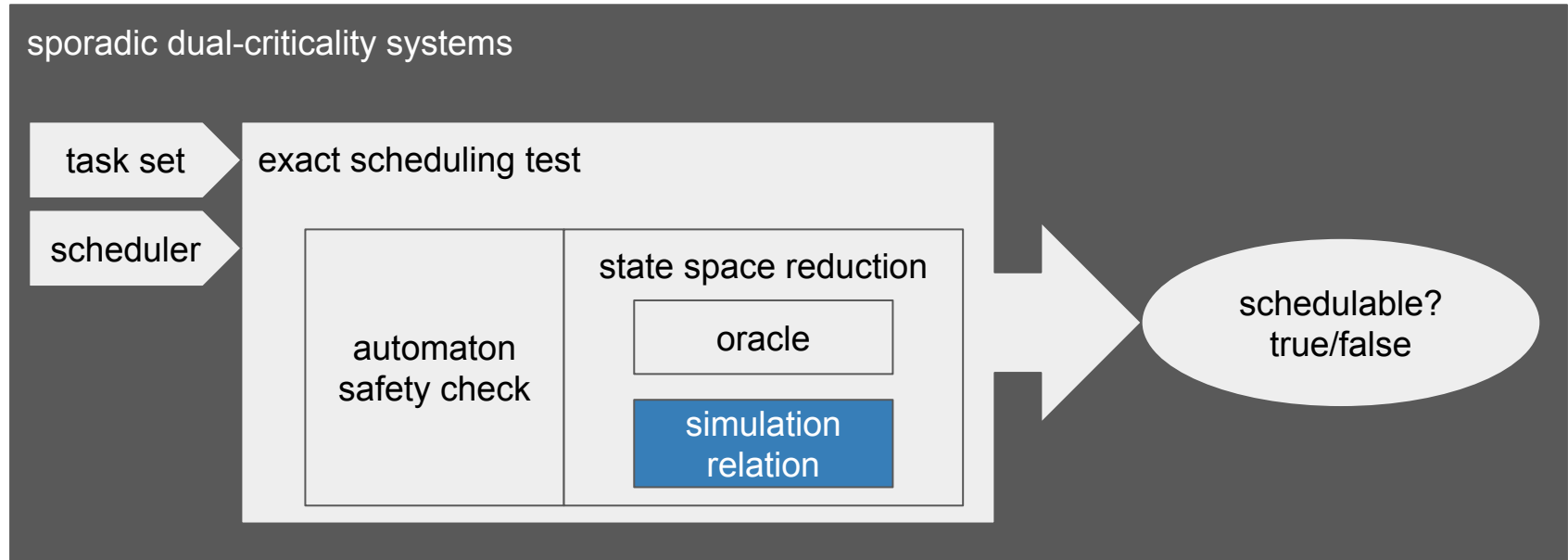
## Visited state reduction (median)



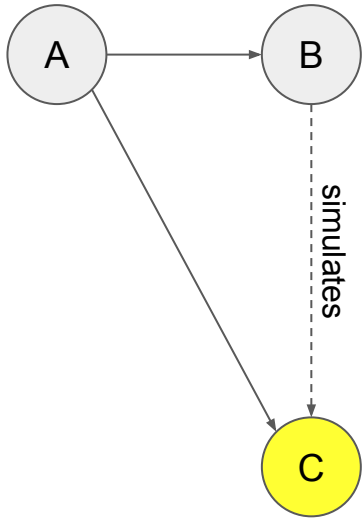
# Objective of the work



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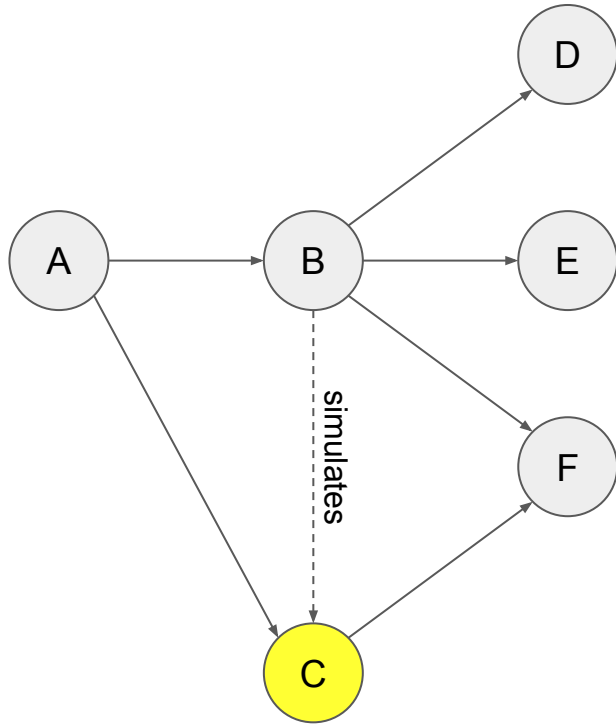
A simulation relation is a relationship between states



state B simulates state C



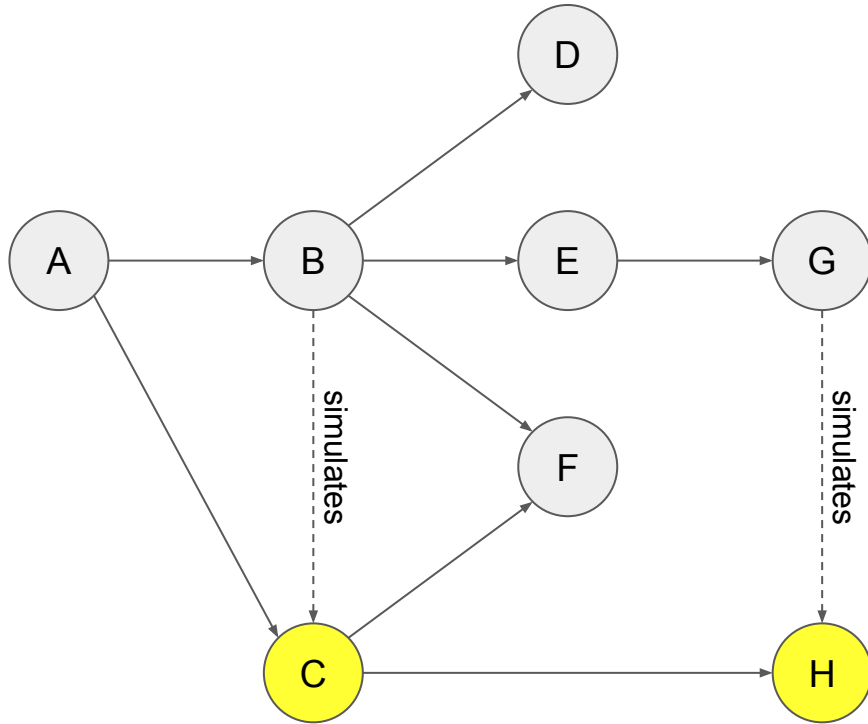
# Implication of the simulation relation



Simulation implies:

1. Can reach all states

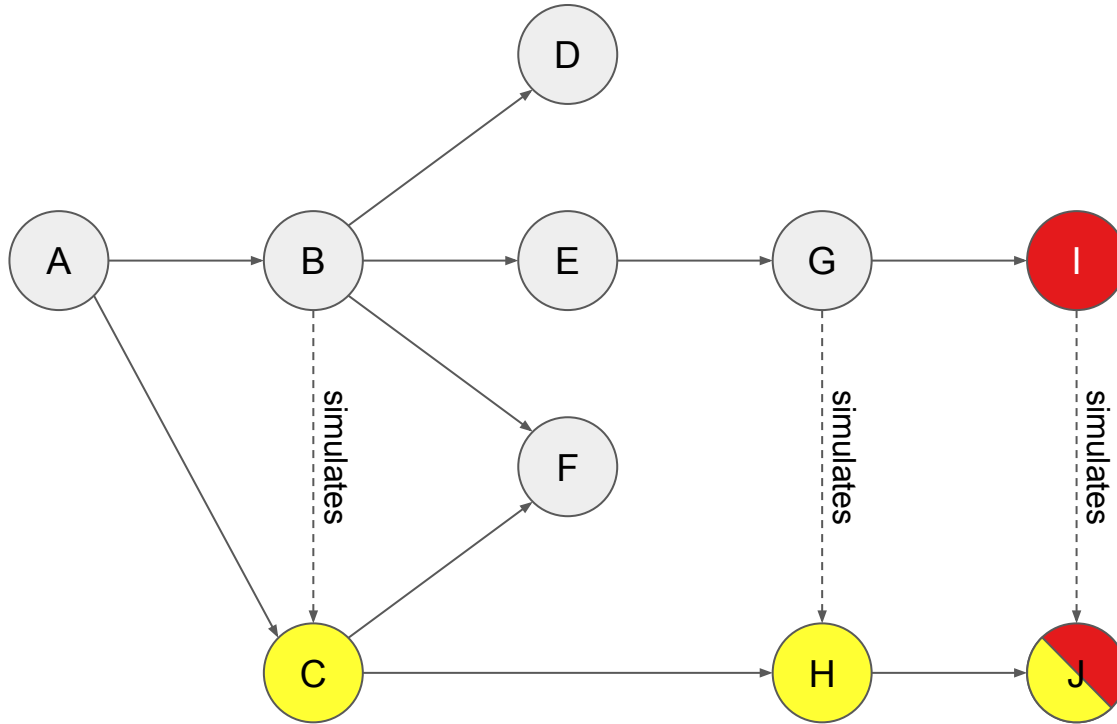
# Implication of the simulation relation



Simulation implies:

1. Can reach all states
2. Or a state simulating it

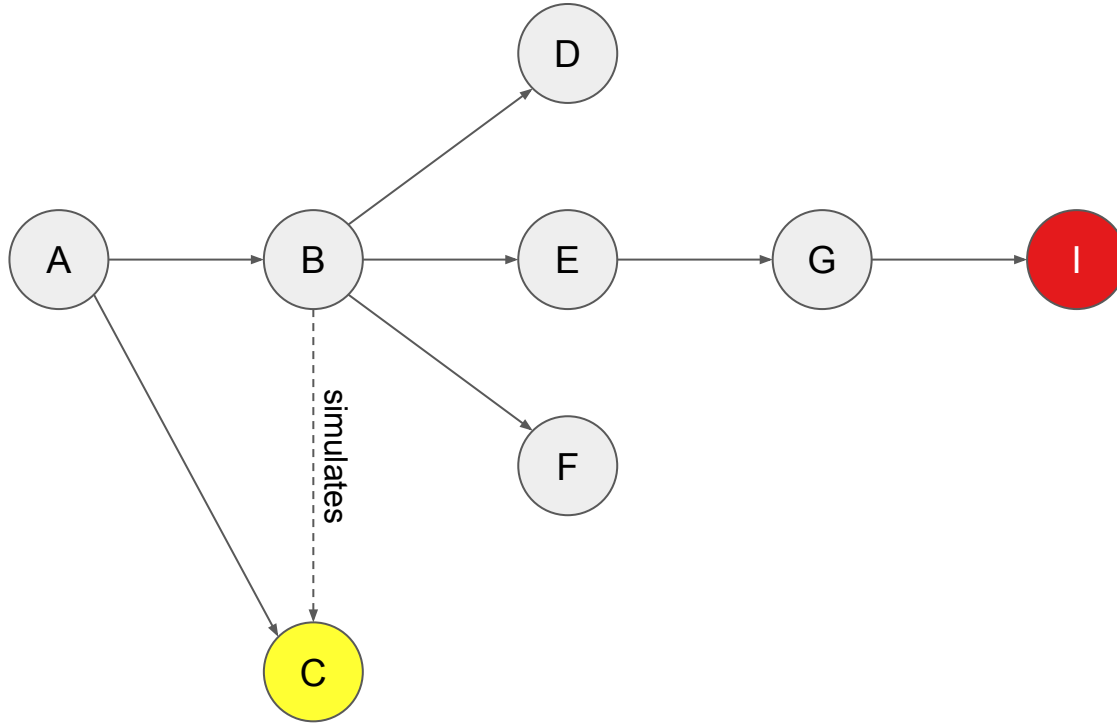
# Implication of the simulation relation



Simulation implies:

1. Can reach all states
2. Or a state simulating it
3. Simulated fail states remains fail states

# Do not explore simulated states and reduce state space



Simulation implies:

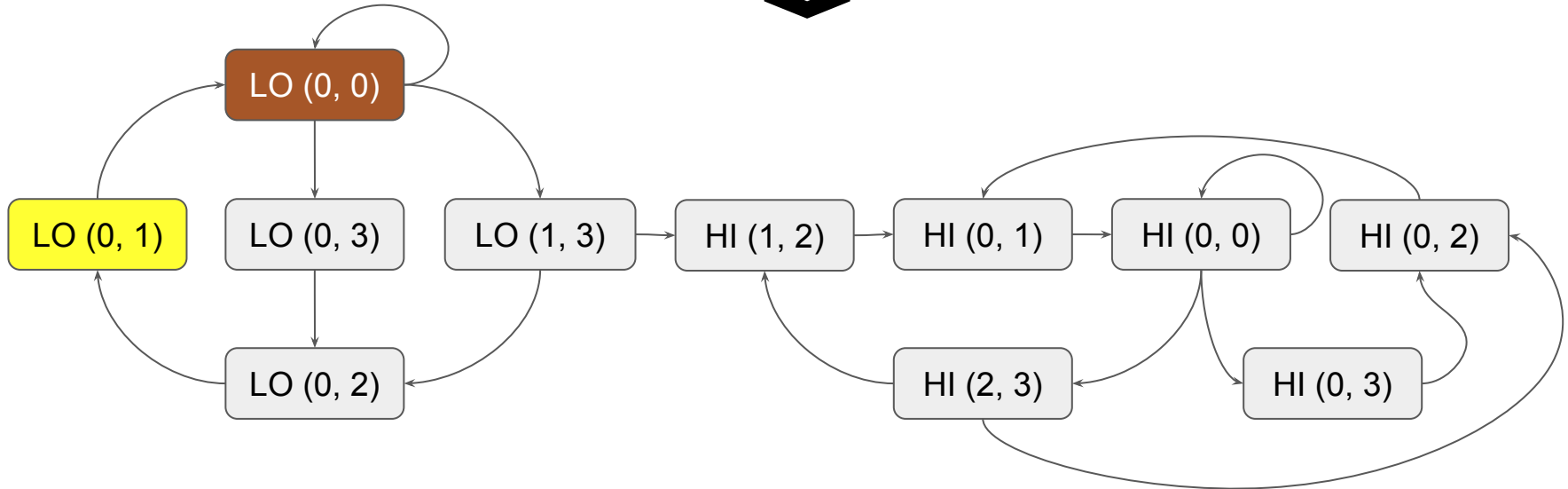
1. Can reach all states
2. Or a state simulating it
3. Simulated fail states remains fail states

# Some states are simulated by the idle state

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



EDF-VD

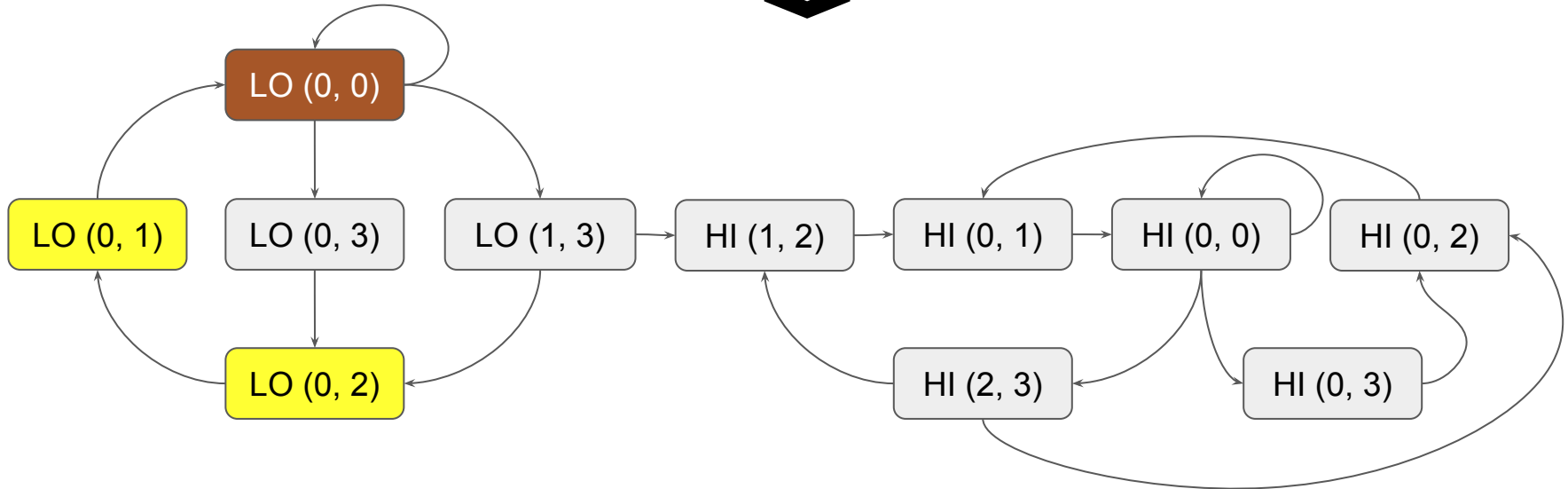


# Some states are simulated by the idle state

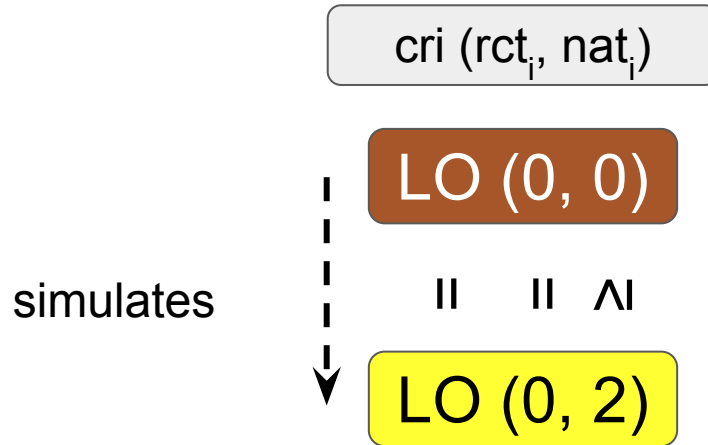
	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



EDF-VD

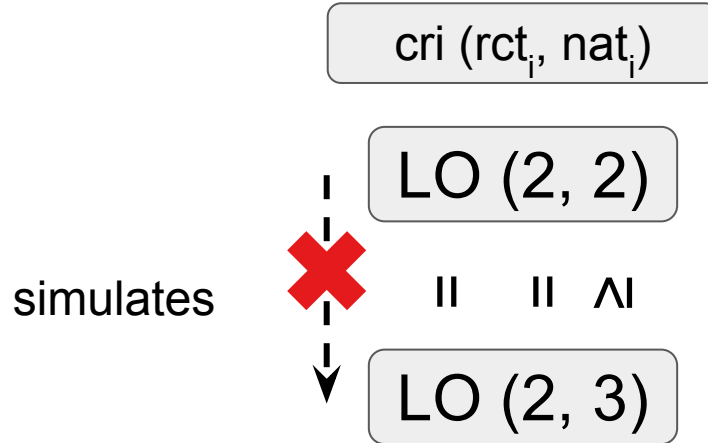


# Idle tasks simulation



- Same criticality
- Task has no active jobs
- Next arrival time is sooner

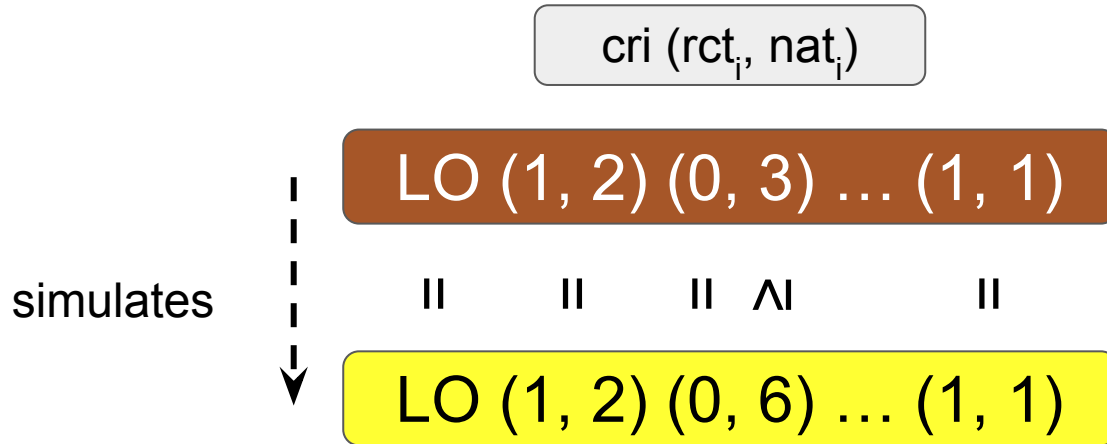
# Idle tasks simulation



- Same criticality
- Active tasks are identical



# Idle tasks simulation



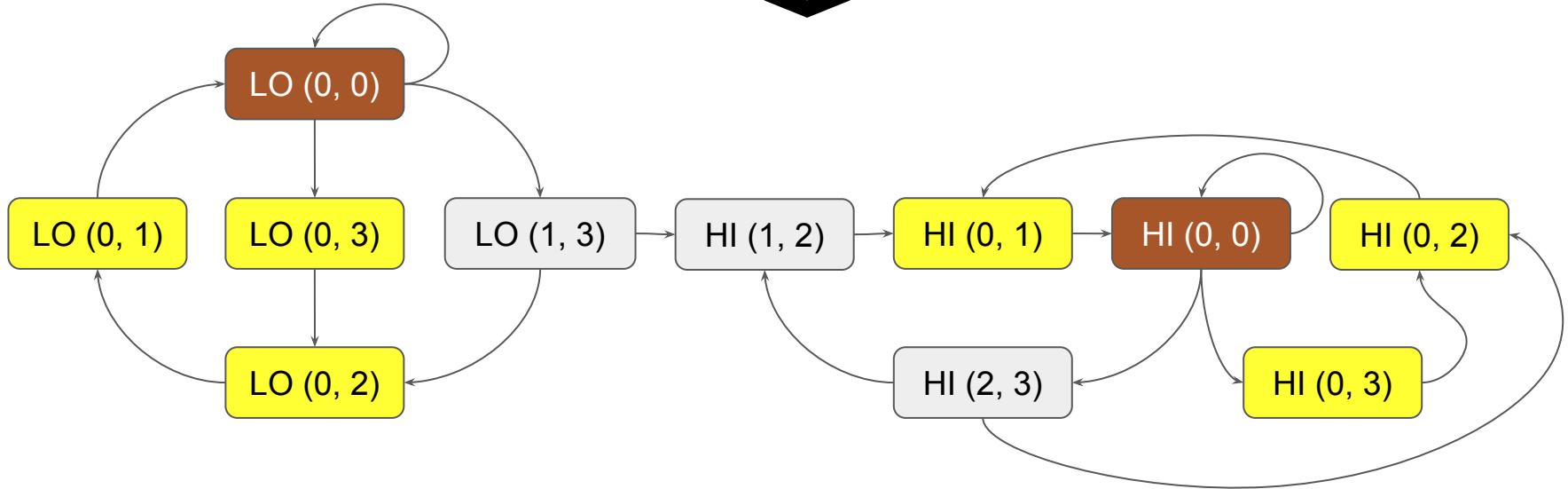
- Same criticality
- Active tasks are identical
- Idle task can release earlier

# Simulated states can be omitted from the exploration

	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI

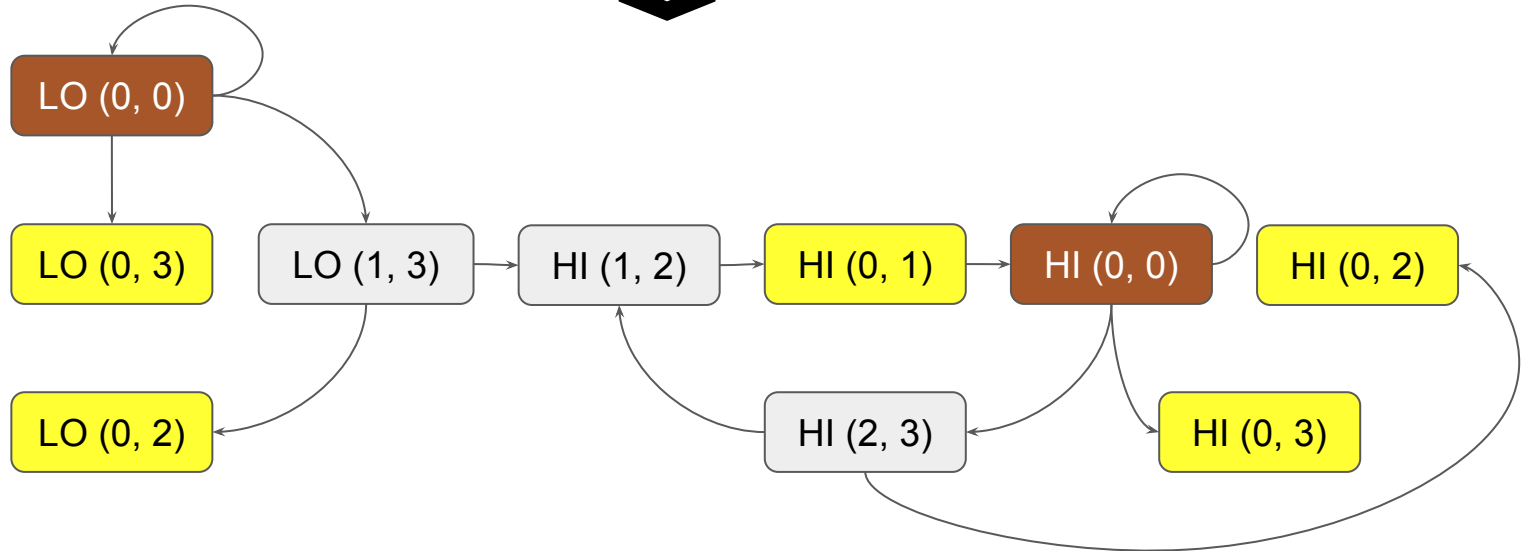


EDF-VD

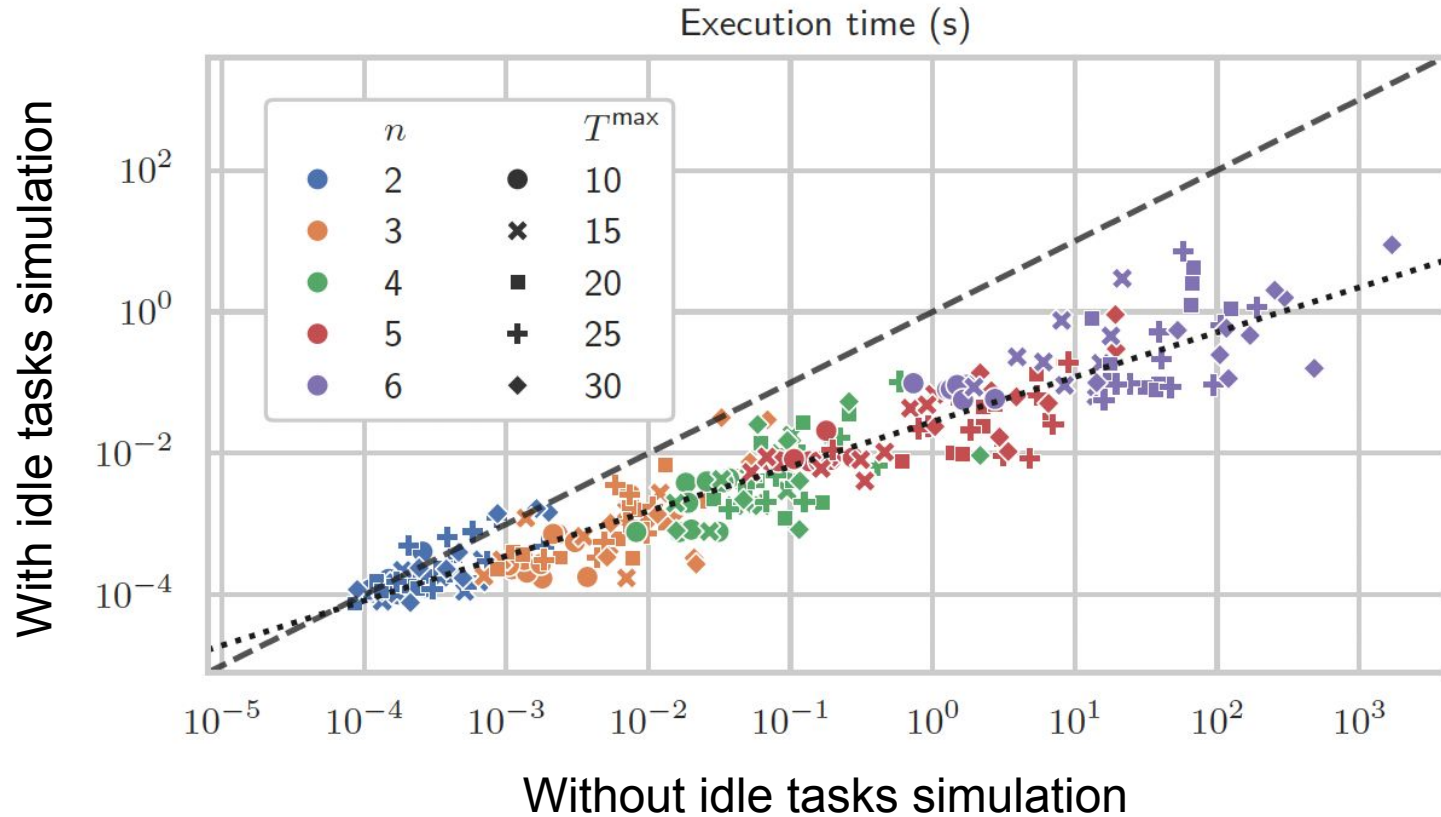


# Simulated states can be omitted from the exploration

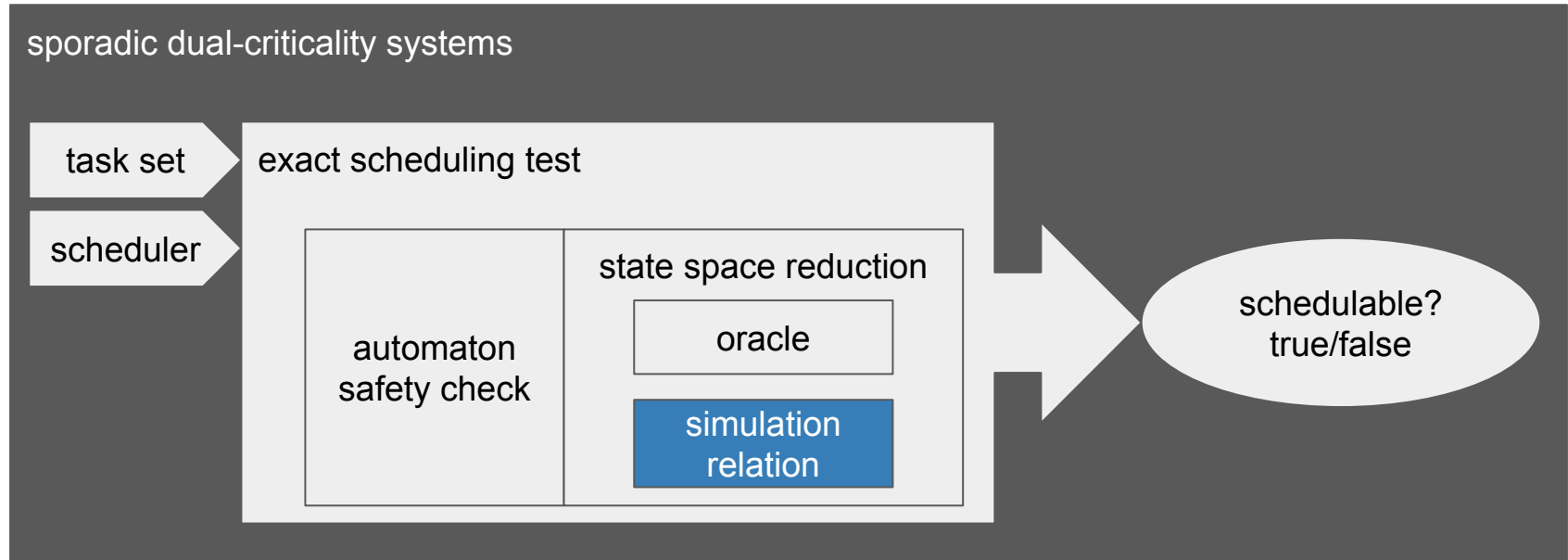
	C(LO)	C(HI)	D	T	L
$\tau_1$	2	3	3	4	HI



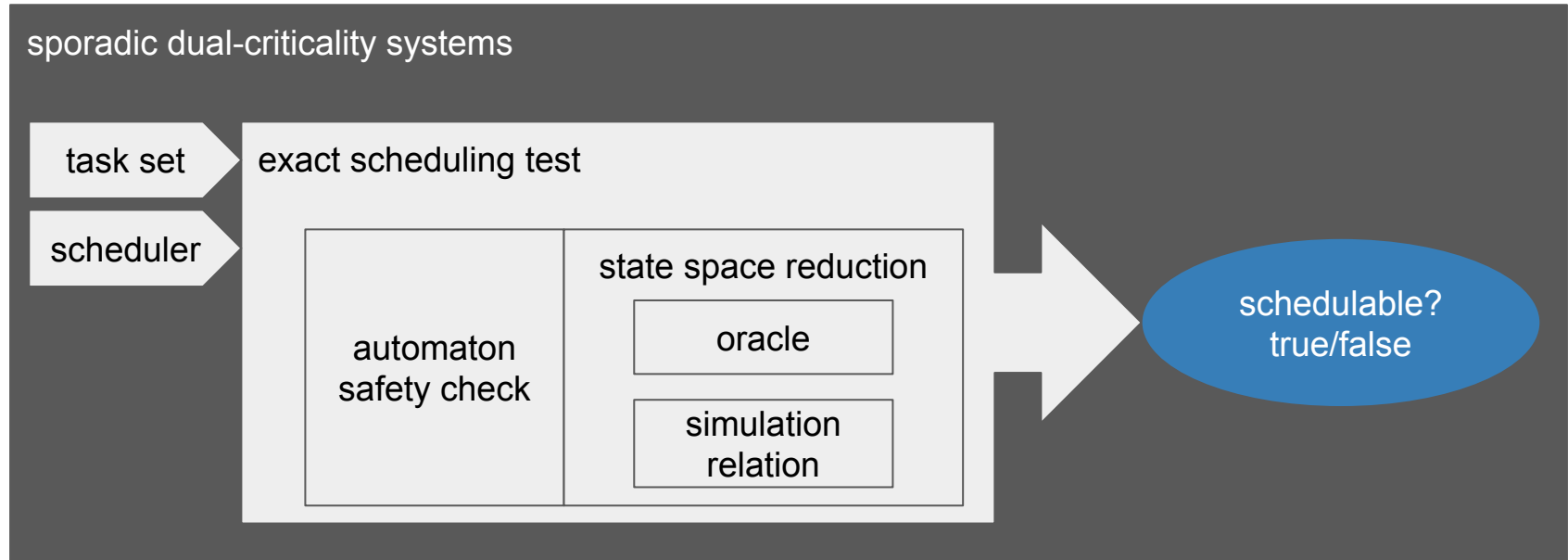
# Idle tasks simulation impact on exploration duration



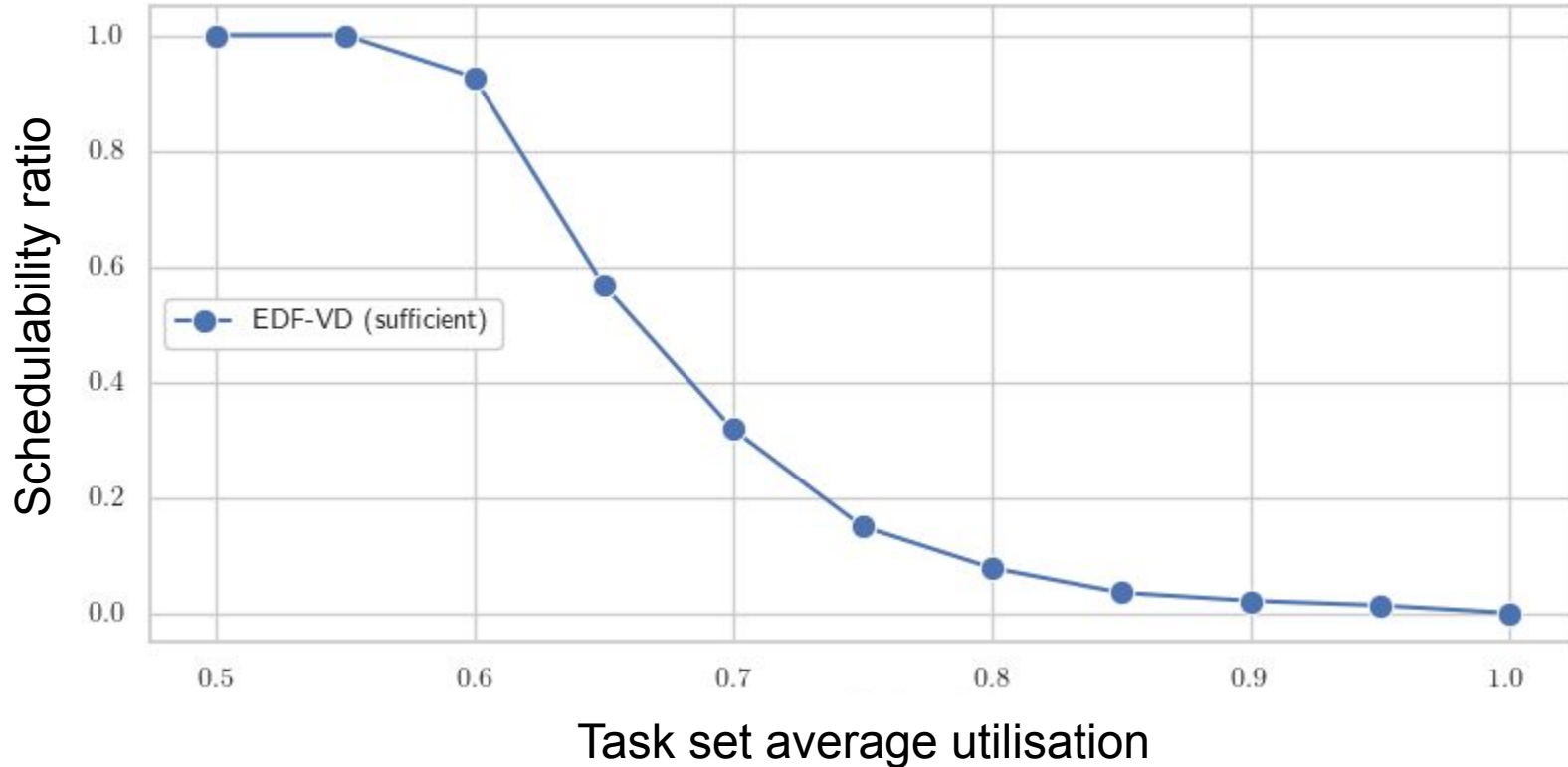
# Objective of the work



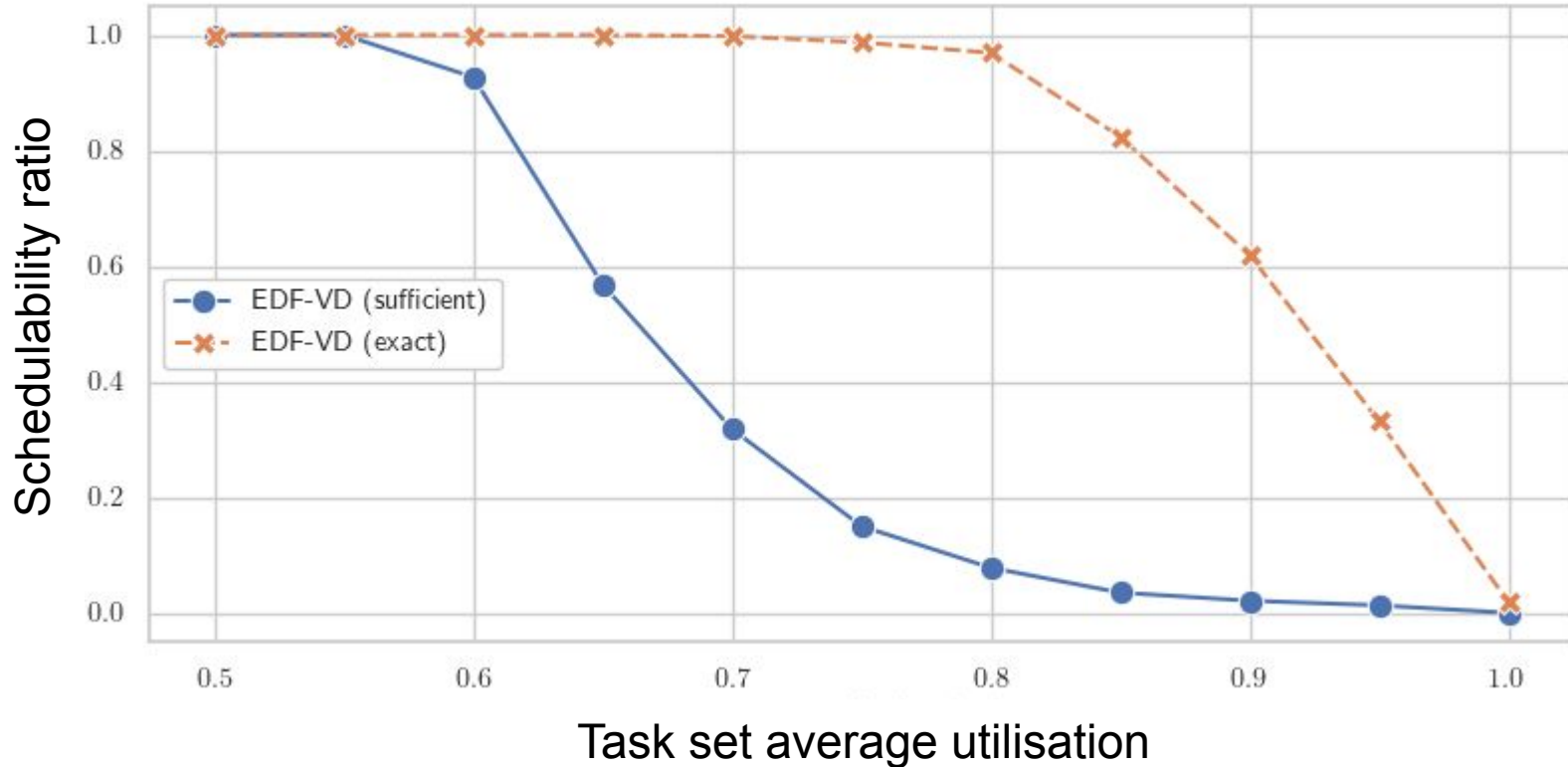
# Objective of the work



# Assessing the schedulability for varying utilisation

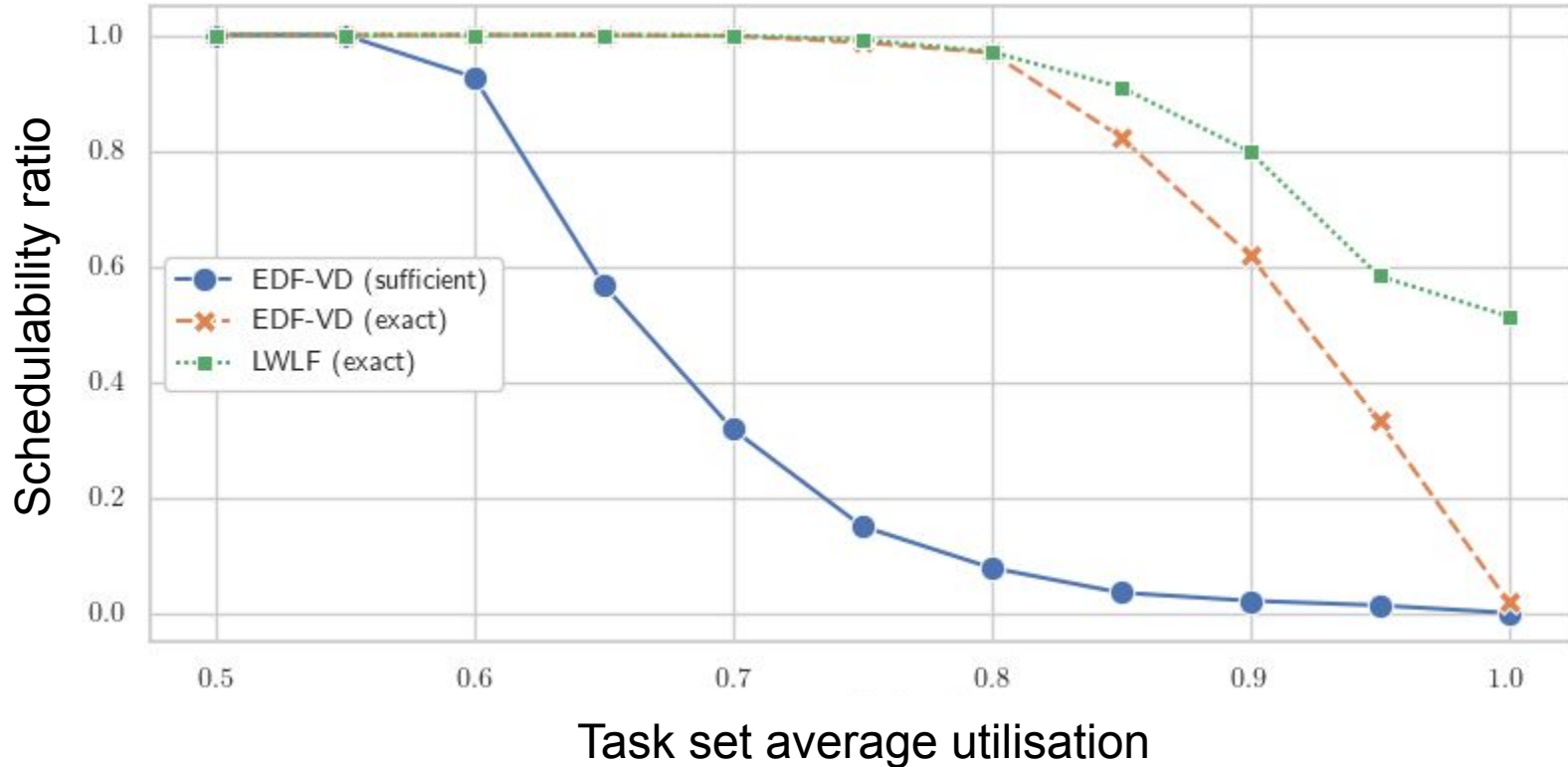


# Quantifying the pessimism of sufficient tests

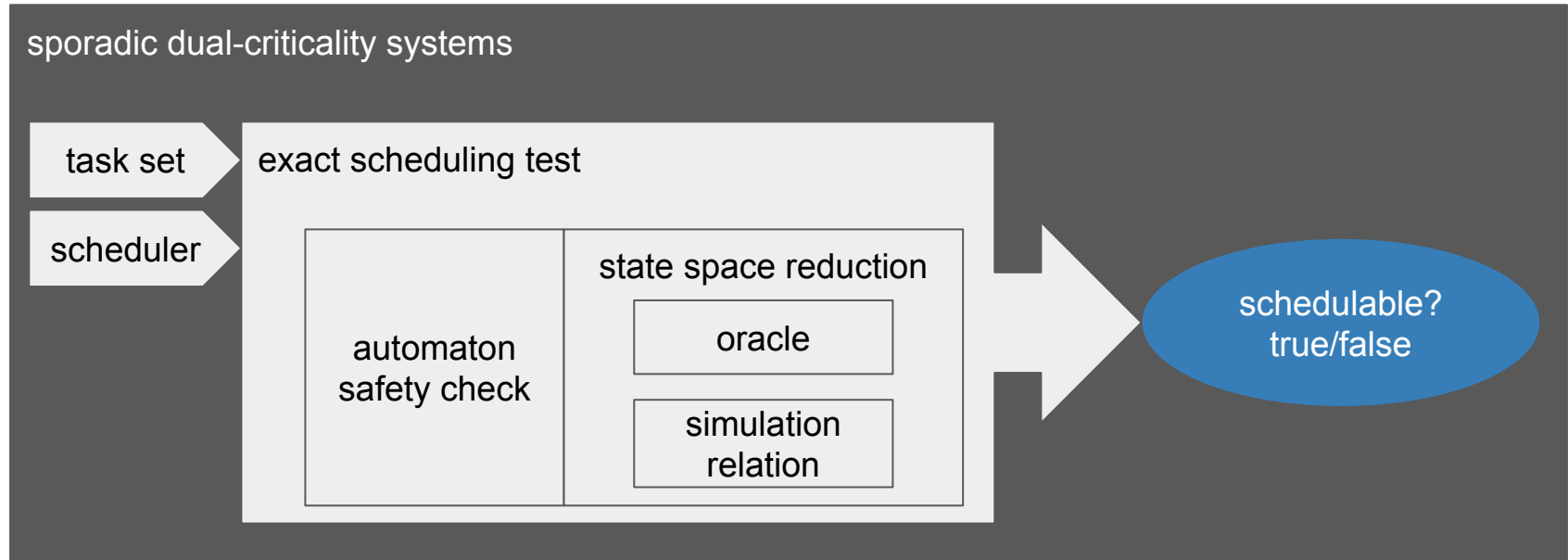




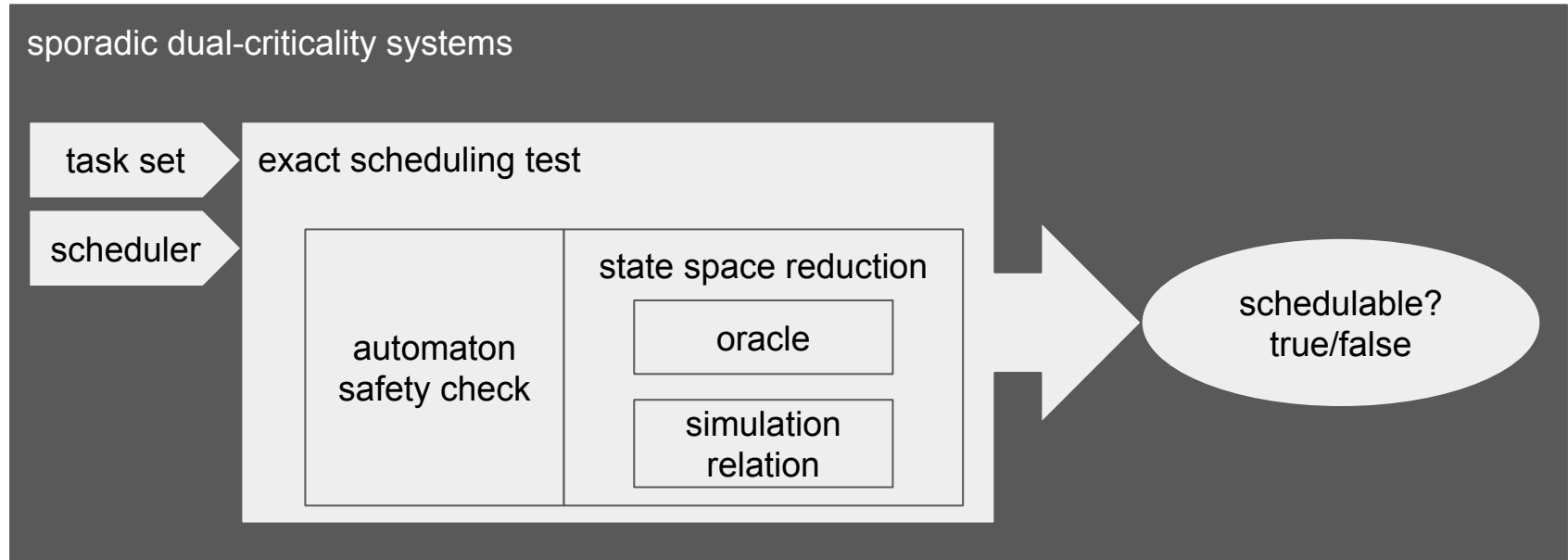
# Empirical scheduler evaluation



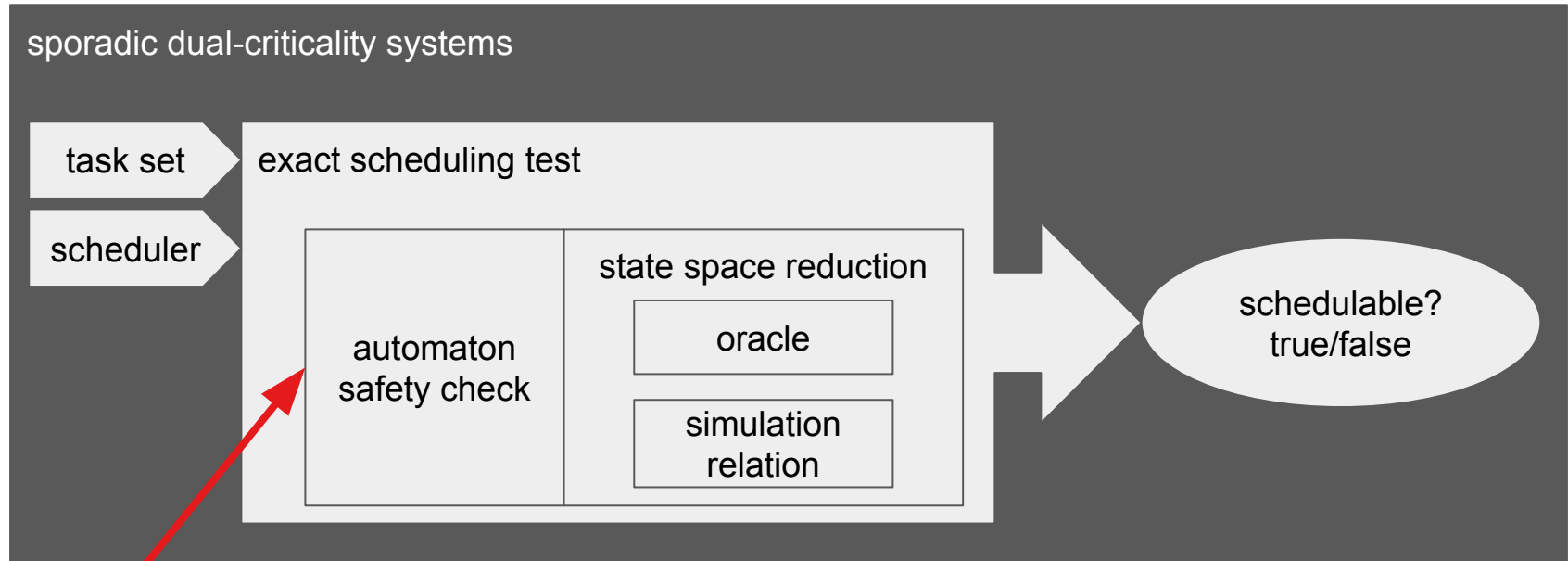
# Objective of the work



# Objective of the work



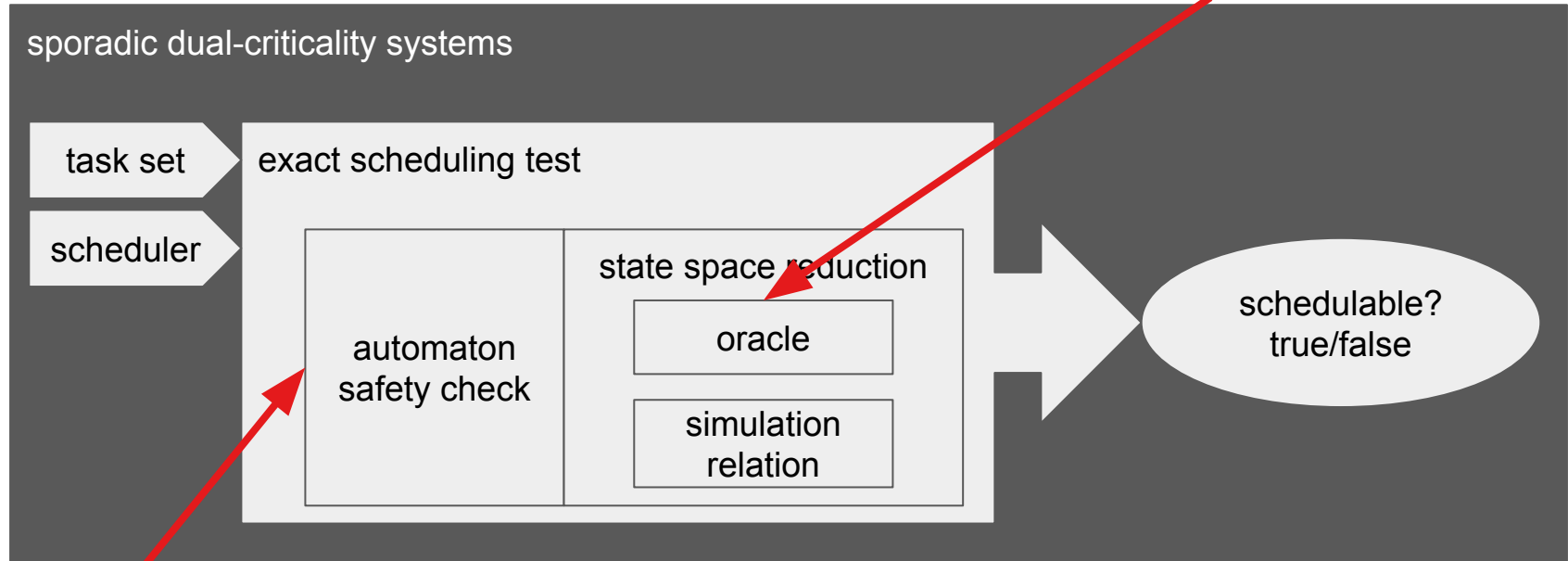
# Objective of the work



**Explicit model  
formal semantics**

# Objective of the work

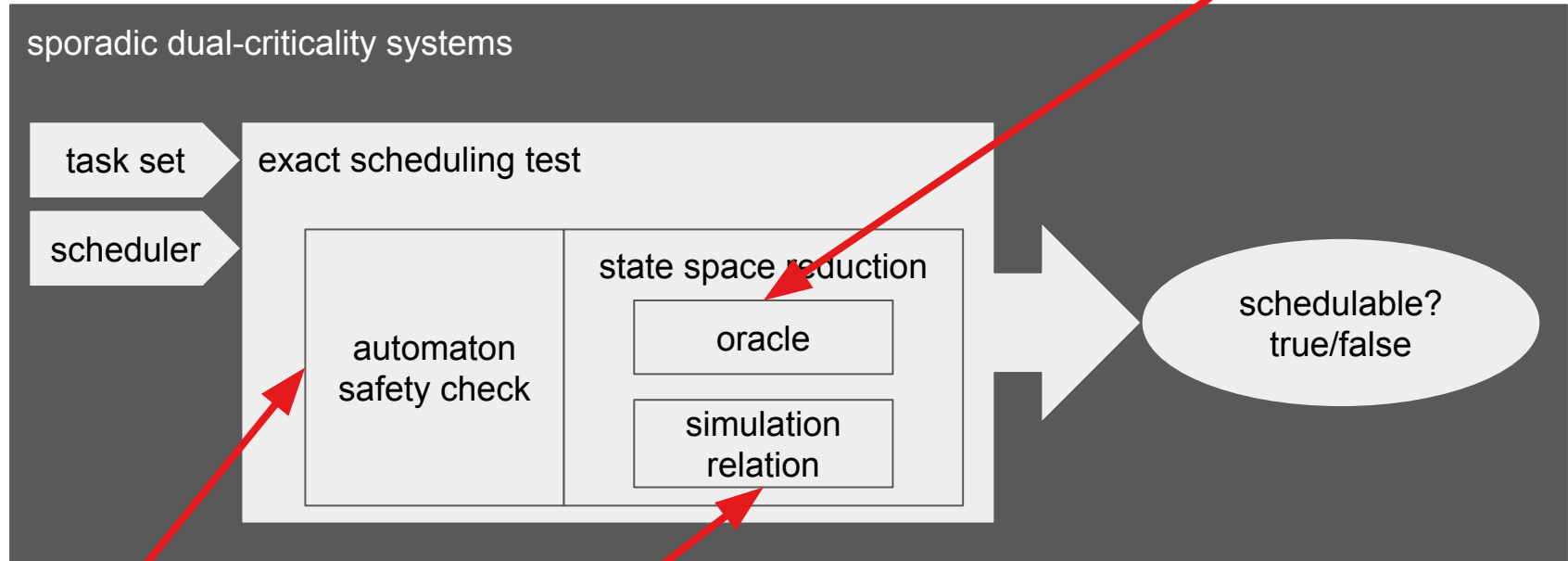
**Real time community knowledge**



**Explicit model  
formal semantics**

# Objective of the work

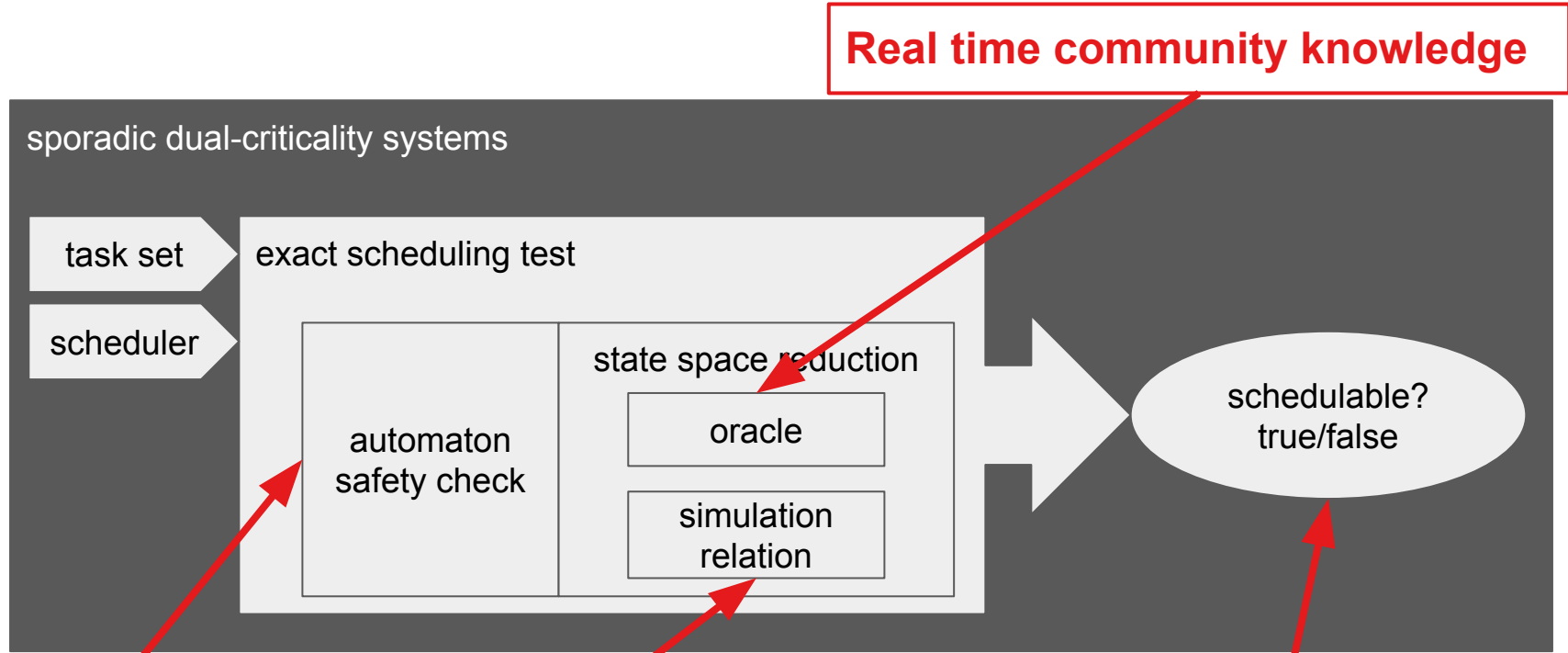
**Real time community knowledge**



**Explicit model formal semantics**

**Formal verification techniques**

# Objective of the work



**Real time community knowledge**

**Explicit model formal semantics**

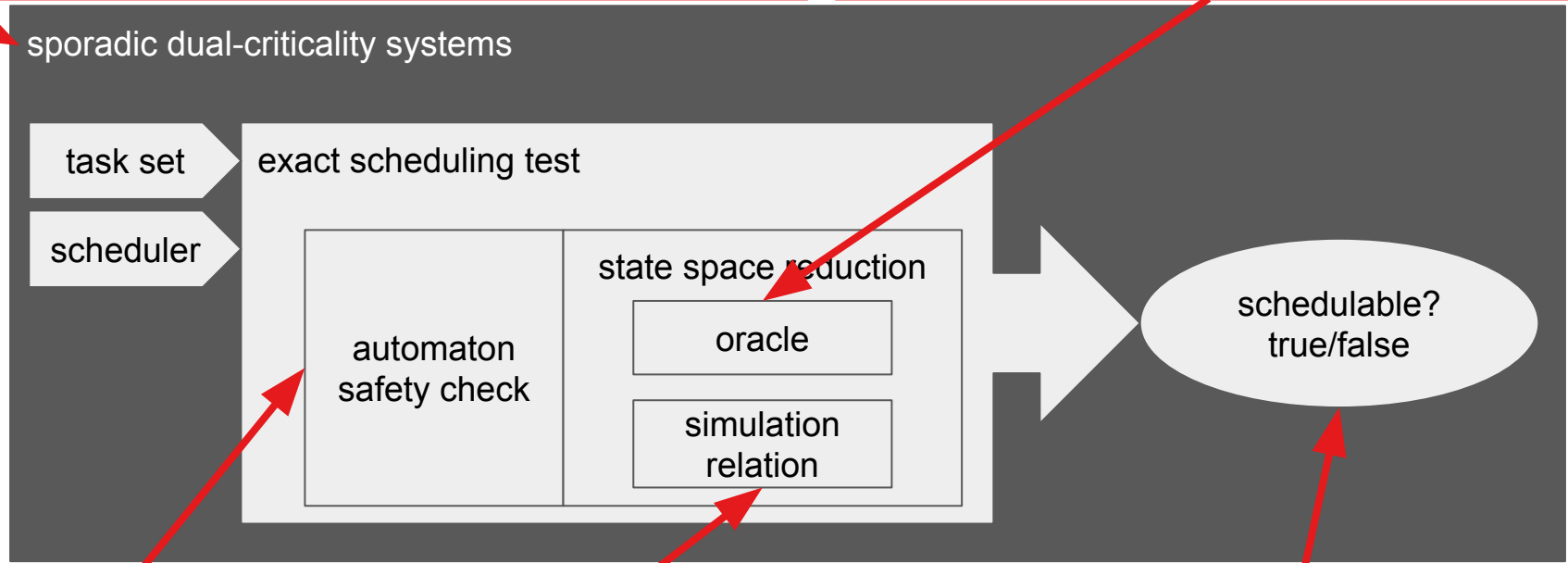
**Formal verification techniques**

**No pessimism and scheduler exploration**

# Objective of the work

**Framework applicable to other models**

**Real time community knowledge**



**Explicit model formal semantics**

**Formal verification techniques**

**No pessimism and scheduler exploration**



# Future work

## **Model extension**

- multi-processor
- varying CPU speed
- preemption delay considerations
- arbitrary deadlines

## **Enhanced scalability**

- new oracles
- new simulation relations
- implement meta transitions

## **Framework extension**

- Feasibility as game on the automaton

# Questions & voting

## Model extension

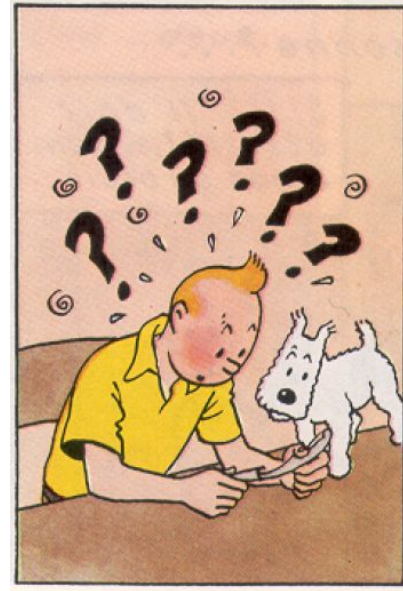
- multi-processor
- varying CPU speed
- preemption delay considerations
- arbitrary deadlines

## Enhanced scalability

- new oracles
- new simulation relations
- implement meta transitions

## Framework extension

- Feasibility as game on the automaton



67 10 96 5

# Backup

# Related work

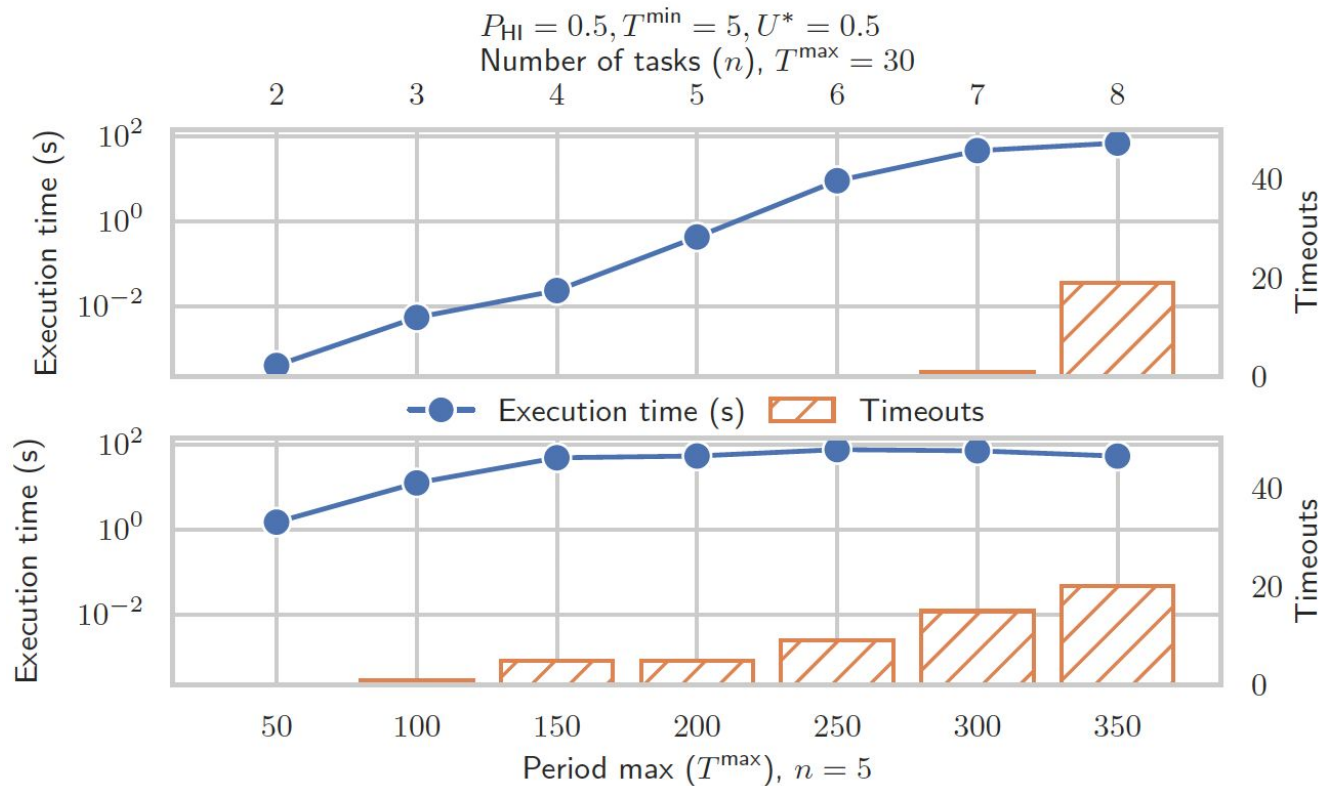
**Table 1:** Comparing with the related work.

	[5]	[4]	[29]	Us
Criticality	Single	Dual	Single	Dual
Priority classes	Any	FTP	Any	Any
Pruning rules	✗	✓	✗	✗
Antichains	✗	✗	✓	✓
Oracles	✗	✗	✗	✓
Multi-processor	✓	✗	✓	✗

# Related work (extended, tentative)

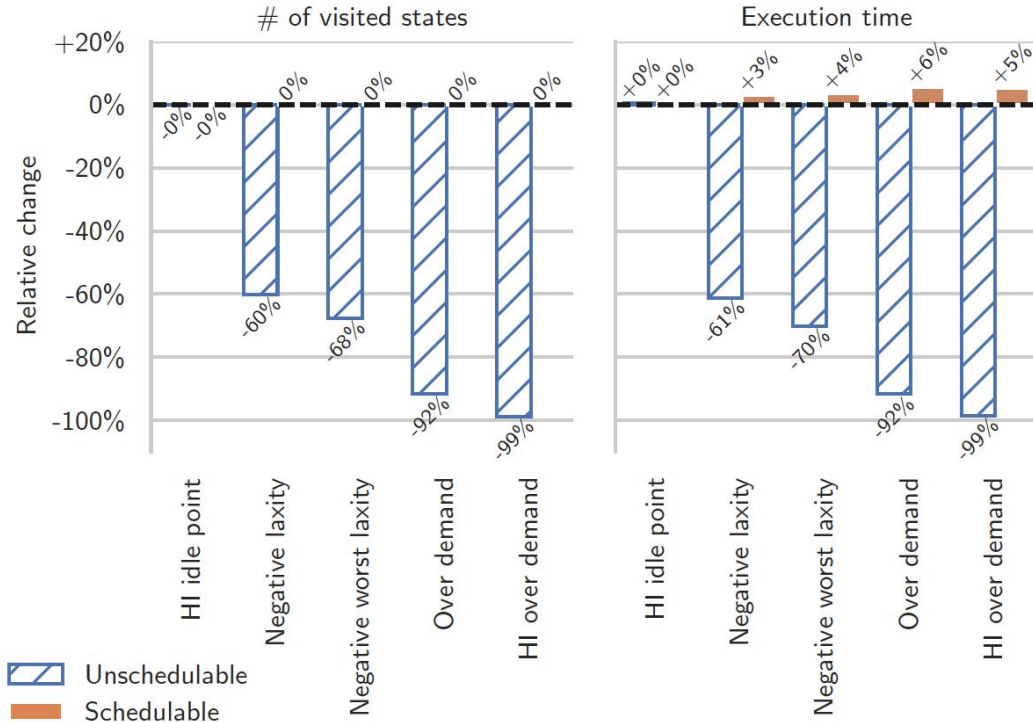
		Baker	Lindstrom	Nasri	Asyaban	Yalcinkaya	Abdeddaïm	Ranjha	Nasri	Gohari	This paper
Category		2007	2011	2017	2018	2019	2020	2023	2024	2024	2024
Framework	FSM	x	x		x						x
	Timed automata					x	x				
	SAG			x				x	x	x	
Execution time	Varying execution time			x		x		x	x	x	
	Dual-crit				x		x				x
Scheduler	FTP	x	x	x	x	x	x	x	x	x	x
	FJP	x	x	x			x	x	x	x	x
	DP	x	x								x
	Preemptive	x	x		x		x		x	x	x
Multi proc	Multi proc	x	x	x		x		x	x	x	
Release model	Release jitter			x		x		x	x	x	
	Periodic task		x	x		x	x		x	x	
	Sporadic task	x	x		x	x	x				x
Constraints	DAG							x			
	Self-suspending					x					
Optimisation	Pruning rules				x			x			
	Partial-order							x		x	
	Antichain		x								x
	Oracles										x

# Scalability experiment



# Extended oracle experiment

$$P_{HI} = 0.5, T^{\min} = 5, T^{\max} = 30, n = 5, U^* \in [0.8; 1; 0.01]$$



# Statistics on the number of visited states

$P_{HI} = 0.5, T^{\min} = 5, T^{\max} = 20, n = 5, U^* \in [0.8; 1; 0.01]$				
Search	BFS		ACBFS	
Oracle	None	HI over demand	None	HI over demand
min	9905	35 (-100%)	668 (-93%)	29 (-100%)
mean	746974	480688 (-36%)	75374 (-90%)	46024 (-94%)
std	999984	819378 (-18%)	126110 (-87%)	107731 (-89%)
median	410063	217928 (-49%)	35888 (-91%)	15459 (-96%)
max	11875126	11875126 (-0%)	2687577 (-77%)	2687577 (-77%)

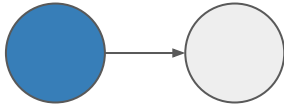


# Safety check in an automaton

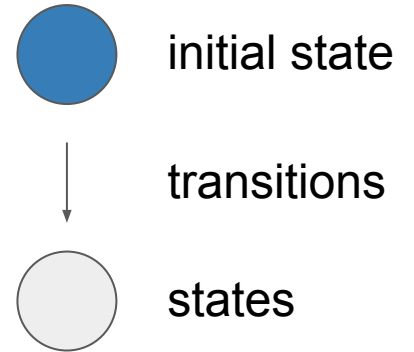
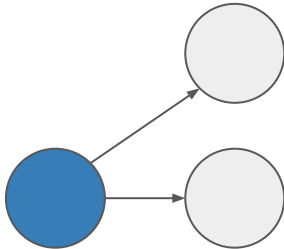


# Safety check in an automaton

automaton #1

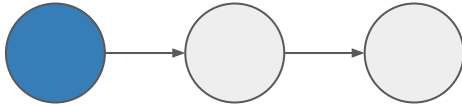


automaton #2

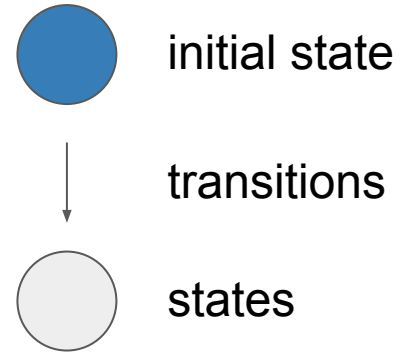
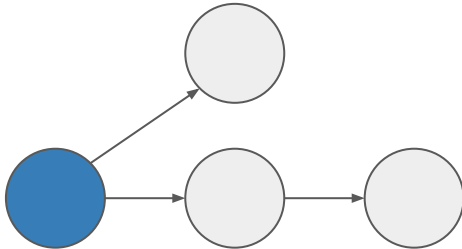


# Safety check in an automaton

automaton #1

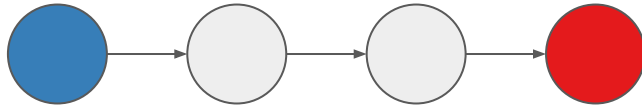


automaton #2



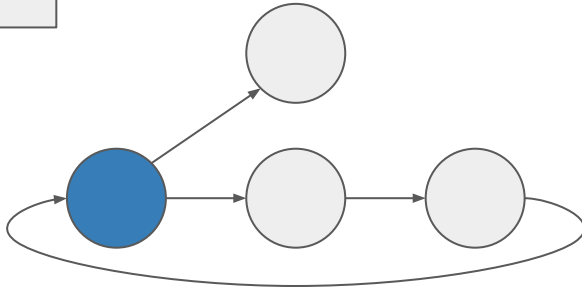
# Safety check in an automaton

automaton #1

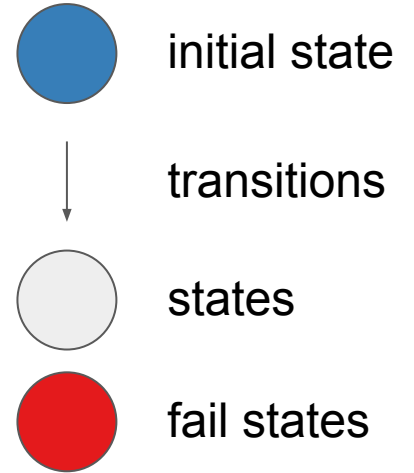


**Unsafe**

automaton #2



**Safe**



# Scheduling test to automaton safety

